

# Lecture 19

CSE 331

Oct 10, 2016

Quiz starts at 1pm  
and ends at 1:10pm

In 1(b), IGNORE time taken to initialize data structures

Lecture starts  
at 1:15pm

# Pitch grading half done

## CSE 331 Mini project choices

Fall 2016

Please check the table below before submitting your mini project pitch to make sure your case study is not being used by another group. Case studies are assigned on a first come first serve basis.

Group	Societal Aspect	Case Studies
Anand Balakrishnan, Vikram Garu and Veronica Ng	Cryptography	Enigma <a href="#">🔗</a> Public key cryptography <a href="#">🔗</a>
Hank Lin, Michael Tobio and Miaomiao Zhang		
Devashish Agarwal, Jacob Fijas and Kevin Rathbun	Social Media in Politics	DeepDrumof <a href="#">🔗</a> Facebook determines your political leanings <a href="#">🔗</a>
Sravanika Doddi, Anne Izydorczak and Simran Singh	Women	Google autocomplete stereotypes <a href="#">🔗</a> Gender bias in job ads <a href="#">🔗</a>

Please wait for grading rubric before asking grading questions

# Interval Scheduling Problem

**Input:**  $n$  intervals  $[s(i), f(i))$  for  $1 \leq i \leq n$

**Output:** A *schedule*  $S$  of the  $n$  intervals

No two intervals in  $S$  conflict

$|S|$  is maximized

# Analyzing the algorithm

$R$ : set of requests

Set  $S$  to be the empty set

While  $R$  is not empty

    Choose  $i$  in  $R$  with the earliest finish time


    Add  $i$  to  $S$

    Remove all requests that conflict with  $i$  from  $R$

Return  $S^* = S$



$S^*$  has no conflicts



$S^*$  is an optimal solution

# Greedy “stays ahead”





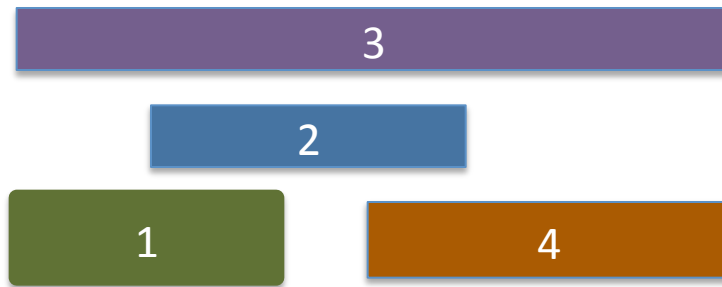
# Today's agenda

Prove the correctness

Analyze run-time of the greedy algorithm

# Algorithm implementation

Go through the intervals in order of their finish time



Check if  $s[i] < f(1)$

with 1:

In general, if  $j$ th interval is the last one chosen

Pick smallest  $i > j$  such that  $s[i] \geq f(j)$  . . .

$O(n \log n)$  run  
time

# The final algo

$O(n \log n)$  time sort intervals such that  $f(i) \leq f(i+1)$

$O(n)$  time build array  $s[1..n]$  s.t.  $s[i]$  = start time for  $i$

Add 1 to  $A$  and set  $f = f(1)$

For  $i = 2 .. n$

    If  $s[i] \geq f$

        Add  $i$  to  $A$

        Set  $f = f(i)$

Return  $A^* = A$

# Reading Assignment

Sec 4.1of [KT]



# Questions?



# The “real” end of Semester blues

There are deadlines and durations of tasks



Write up a term paper

Party!

Exam study

331 HW

Project

Monday

Tuesday

Wednesday

Thursday

Friday

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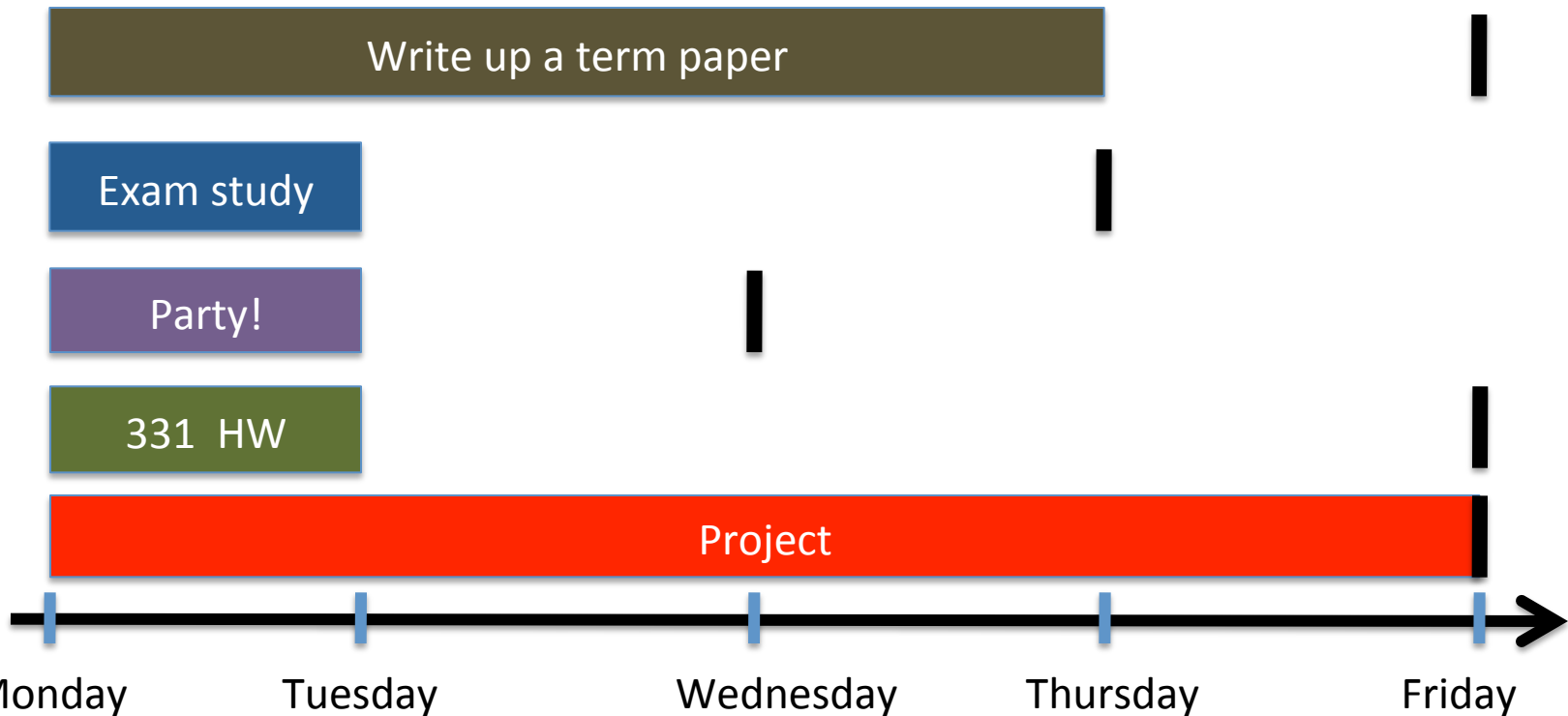
Monday

Tuesday

Wednesday

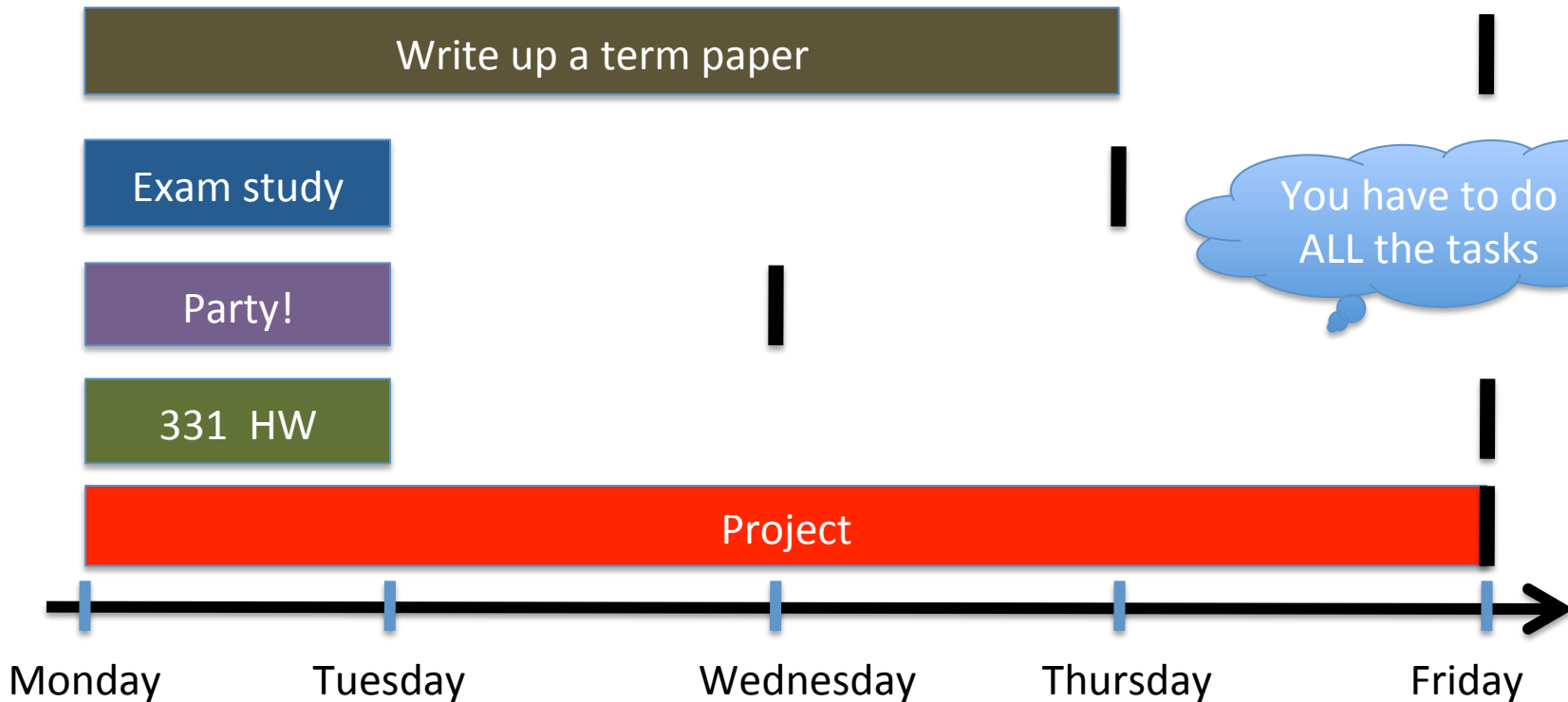
Thursday

Friday



# The algorithmic task

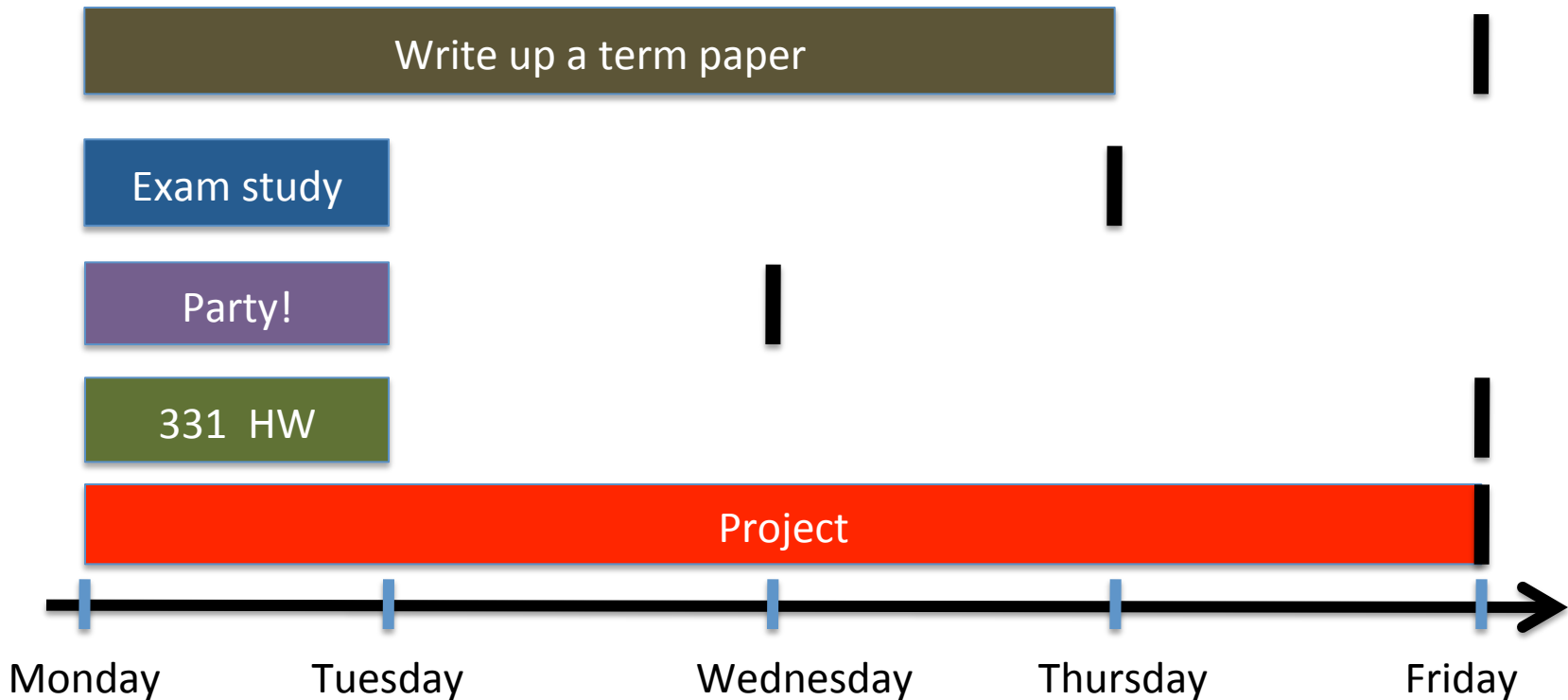
YOU decide when to start each task





# Scheduling to minimize lateness

All the tasks have to be scheduled  
GOAL: minimize maximum lateness



# One possible schedule

All the tasks have to be scheduled  
GOAL: minimize maximum lateness

