

# Lecture 19

CSE 331

Oct 12, 2016

# Mid-term-I Monday

In class

1:00pm-1:50pm sharp

Eight True/False with justification Qs

# Questions?



# Analyzing the algorithm

$R$ : set of requests

Set  $A$  to be the empty set

While  $R$  is not empty

    Choose  $i$  in  $R$  with the earliest finish time

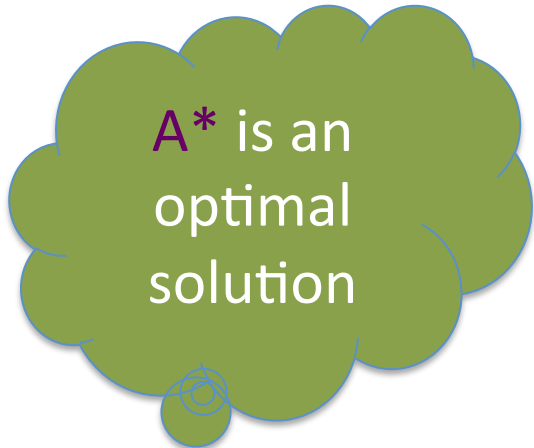
    Add  $i$  to  $A$

    Remove all requests that conflict with  $i$  from  $R$

Return  $A^* = A$



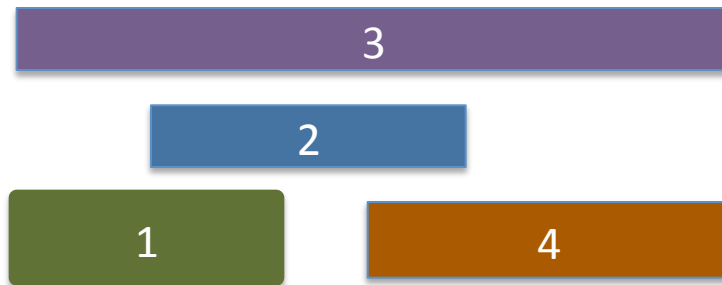
$A^*$  has no conflicts



$A^*$  is an optimal solution

# Algorithm implementation

Go through the intervals in order of their finish time



Check if  $s[i] < f(1)$

with 1:

In general, if  $j$ th interval is the last one chosen

Pick smallest  $i > j$  such that  $s[i] \geq f(j)$  . . .

$O(n \log n)$  run  
time

# The final algo

$O(n \log n)$  time sort intervals such that  $f(i) \leq f(i+1)$

$O(n)$  time build array  $s[1..n]$  s.t.  $s[i]$  = start time for  $i$

Add 1 to  $A$  and set  $f = f(1)$

For  $i = 2 .. n$

    If  $s[i] \geq f$

        Add  $i$  to  $A$

        Set  $f = f(i)$

Return  $A^* = A$

# Reading Assignment

Sec 4.1of [KT]



# Questions?





# The “real” end of Semester blues

There are deadlines and durations of tasks



Write up a term paper

Party!

Exam study

331 HW

Project

Monday

Tuesday

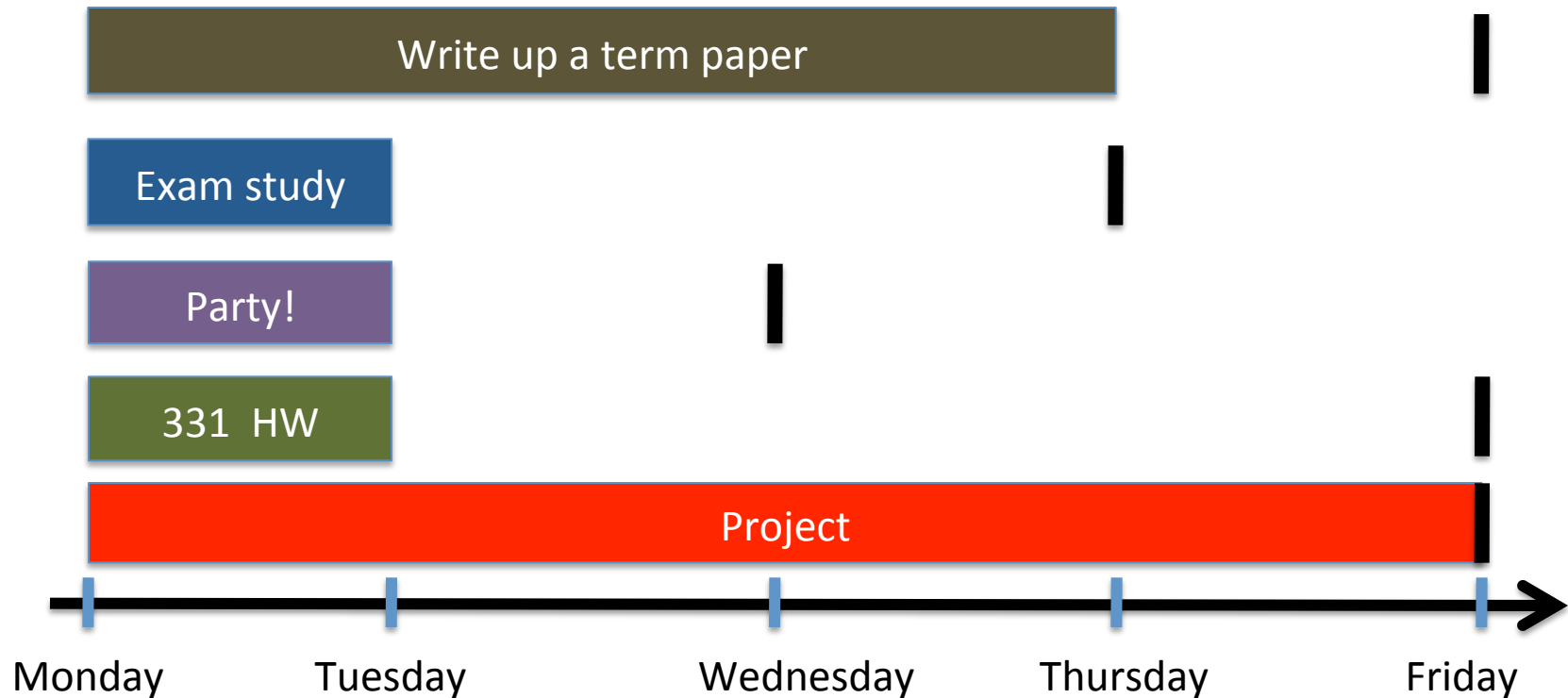
Wednesday

Thursday

Friday

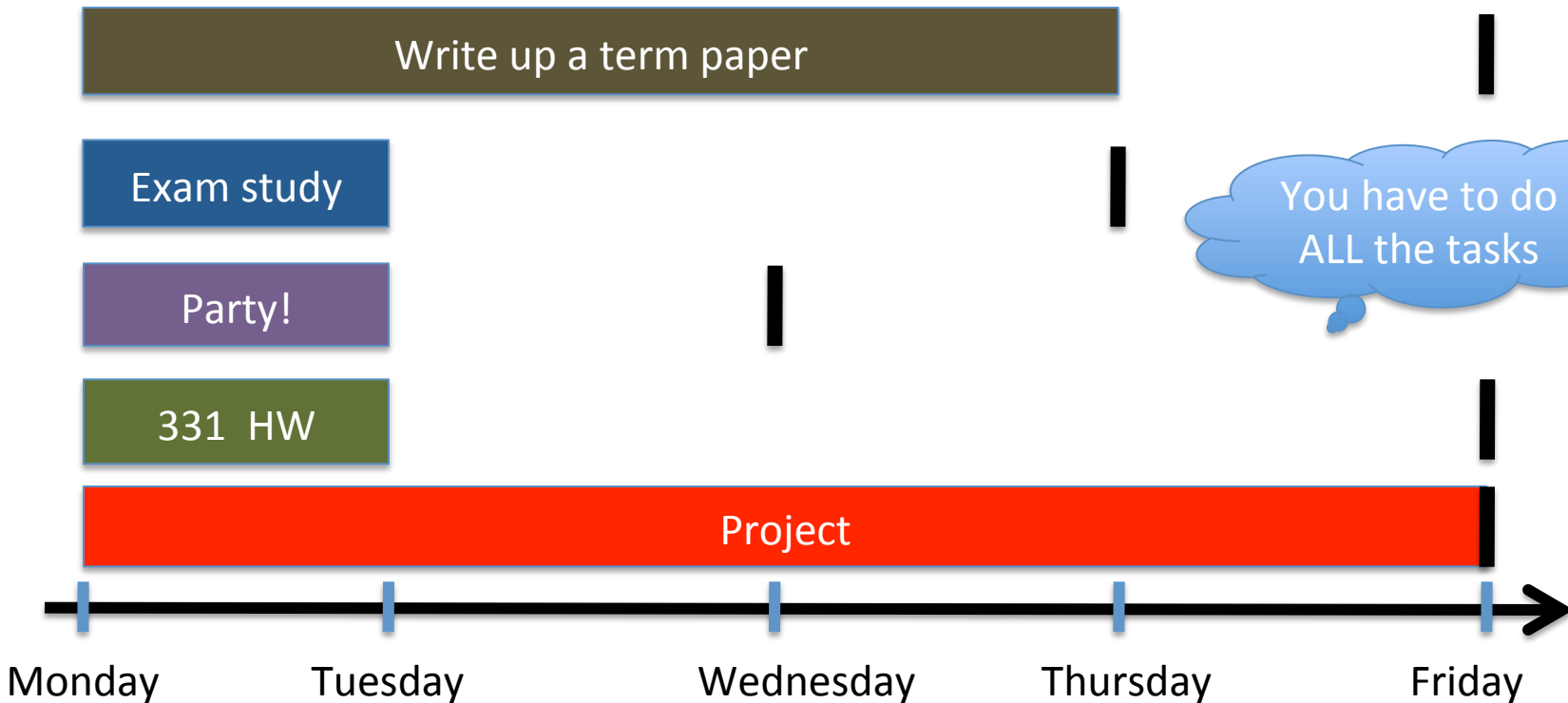
# The “real” end of Semester blues

There are deadlines and durations of tasks



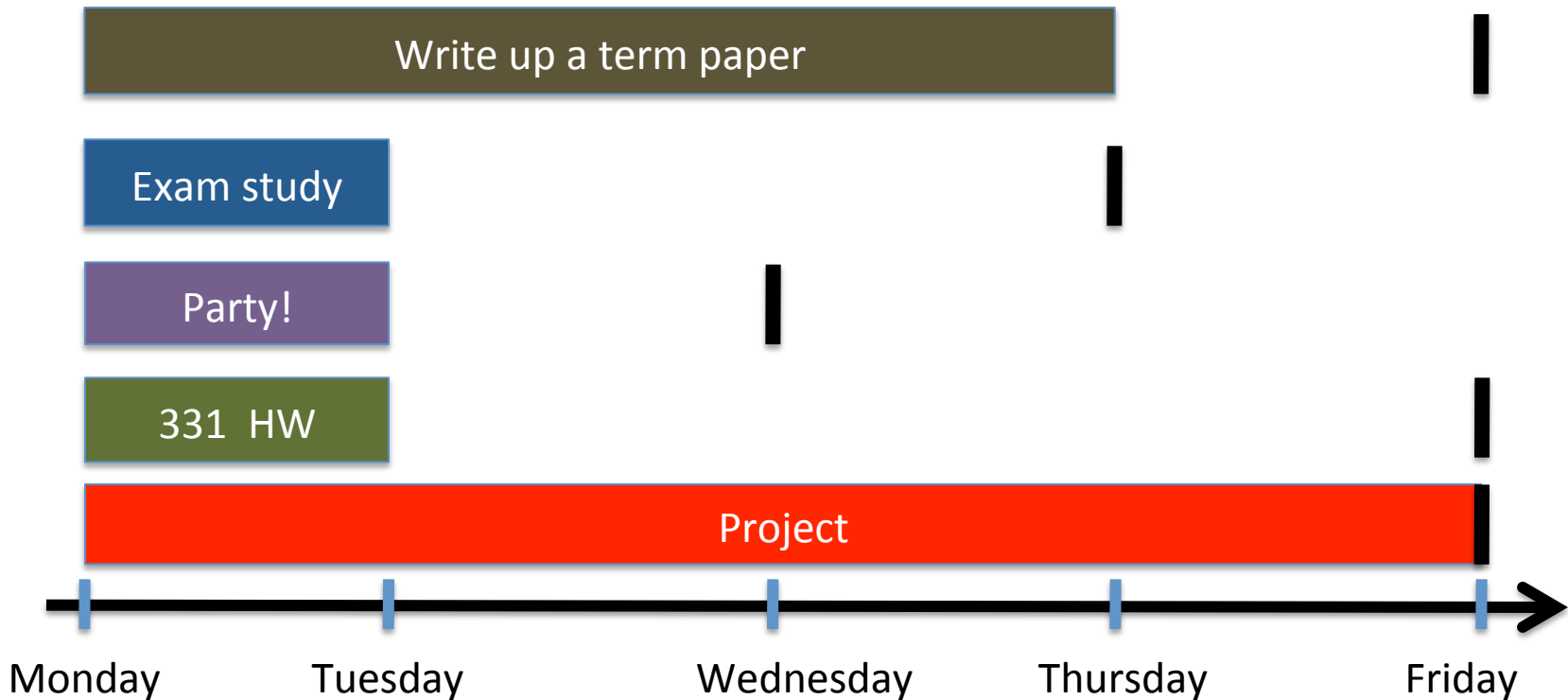
# The algorithmic task

YOU decide when to start each task



# Scheduling to minimize lateness

All the tasks have to be scheduled  
GOAL: minimize maximum lateness



# One possible schedule

All the tasks have to be scheduled  
GOAL: minimize maximum lateness

