## Lecture 35

CSE 331
Nov 28, 2016

## Quiz 2 next Monday

## Quiz 2 in two weeks

A gentle reminder that quiz 2 will be in class on Monday, December 5 from 1-1:10pm.
The first two questions will be T/F without justification (so like two from Q1 on sample final- @735) and the third question will be T/F with justification (so like one from Q2 on sample final- @735).
\#pin
quiz2

## Comments on Feedback

## Comments on feedback

Thanks for everyone who have feedback (@627). Over the course of this week, I will address/respond to some of the feedback (both the quantitative ones and the written comments).

In some cases I will be able to incorporate your comments this year. For others, it might not be but I will at least present you my rationale for for why not.

To being with here is the overall response:

I will pin this post and over the course of the week, I will link to separate posts that had more comments on various sections on the feedback:

1. Homeworks (@759)
2. Lectures (@760)
3. TAs (@761)
4. Exams (@762)
5. Overall comments (@763)

## Official Feedback forms

## Incentive for filling in the course evaluations

You must have received an email about filling the course evaluation forms. I believe this is the link:
https://buffalo.campuslabs.com/courseeval/

Taking a leaf out of Geoff's playbook, here is my offer to incentivize you guys filling in the course evaluation form:

- If at least $\mathbf{8 5 \%}$ o you fill in the course evaluation form, then I will release one T/F (without justification) question on the final exam (which correspor is to Q1 (a): see @735 for the format).
- If at least $95 \%$
you fill in the course evaluation form, then I will release one T/F (without justification) question and one T/F question (corresponding to Q1(a) and Q2(a) respectively: see @735 for the format).

Of course if $<85 \%$ of you fill in the course eval form, then no question gets released. The current percentage of you who have filled in the course evals is $3 \%$ (so a bit to go). I will post weekly updates on the response rate.

Course Evaluations
$\Omega$ Active
11 Days Remaining
Ends 12/8/2016 at 11:59 PM EST

## CS Ed week (Dec 5)

We need volunteers!

with the Department of Computer Science and Engineering at UB:

Children K-12 are invited to:

We need demos!

## KID'S <br> 

Monday, Dec. 5 | Davis Hall, UB


## When to use Dynamic Programming

There are polynomially many sub-problems
OPT(1), ..., OPT(n)


Richard Bellman

Optimal solution can be computed from solutions to sub-problems

$$
\text { OPT }(j)=\max \left\{v_{j}+O P T(p(j)), O P T(j-1)\right\}
$$

There is an ordering among sub-problem that allows for iterative solution

OPT (j) only depends on OPT(j-1), ..., OPT(1)

## Scheduling to min idle cycles

$n$ jobs, ith job takes $w_{i}$ cycles

You have W cycles on the cloud

What is the maximum number of jobs you can schedule?

## Subset sum problem

Input:
$n$ integers $w_{1}, w_{2}, \ldots, w_{n}$
bound W

Output: subset $S$ of $[n]$ such that
(1) sum of $w_{i}$ for all i in $S$ is at most $W$
(2) $w(S)$ is maximized

## Recursive formula

OPT(j, W') = max value out of $\mathrm{W}_{1}, . ., \mathrm{w}_{\mathrm{j}}$ with bound $\mathrm{W}^{\prime}$

If $w_{j}>W^{\prime}$
OPT(j, W') = OPT(j-1, W')
else

$$
\text { OPT(j, W') }=\max \left\{O P T\left(j-1, W^{\prime}\right), w_{j}+O P T\left(j-1, W^{\prime}-w_{j}\right)\right\}
$$

## Today's agenda

Dynamic Program for Subset Sum problem

## Shortest Path Problem

Input: (Directed) Graph $\mathrm{G}=(\mathrm{V}, \mathrm{E})$ and for every edge e has a cost $\mathrm{c}_{\mathrm{e}}$ (can be $<0$ )
$t$ in $V$

Output: Shortest path from every s to $t$


Assume that G has no negative cycle

## May the Bellman force be with you



