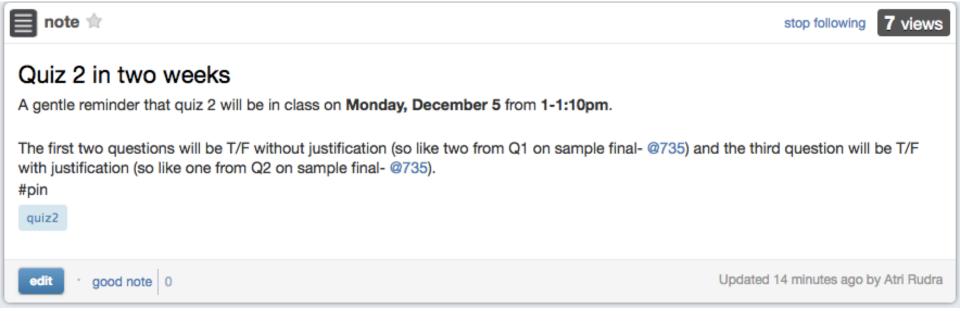
Lecture 35

CSE 331

Nov 28, 2016

Quiz 2 next Monday



Comments on Feedback



stop following

118 views

Actions ▼

Comments on feedback

Thanks for everyone who have feedback (@627). Over the course of this week, I will address/respond to some of the feedback (both the quantitative ones and the written comments).

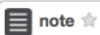
In some cases I will be able to incorporate your comments this year. For others, it might not be but I will at least present you my rationale for for why not.

To being with here is the overall response:

I will pin this post and over the course of the week, I will link to separate posts that had more comments on various sections on the feedback:

- 1. Homeworks (@759)
- 2. Lectures (@760)
- 3. TAs (@761)
- 4. Exams (@762)
- 5. Overall comments (@763)

Official Feedback forms



1 views

Actions ▼

Incentive for filling in the course evaluations

You must have received an email about filling the course evaluation forms. I believe this is the link:

https://buffalo.campuslabs.com/courseeval/

Taking a leaf out of Geoff's playbook, here is my offer to incentivize you guys filling in the course evaluation form:

- If at least 85% of you fill in the course evaluation form, then I will release one T/F (without justification) question on the final exam (which corresponds to Q1(a): see @735 for the format).
- If at least 95% c you fill in the course evaluation form, then I will release one T/F (without justification) question and one T/F (with justification) question (corresponding to Q1(a) and Q2(a) respectively: see @735 for the format).

Of course if < 85% of you fill in the course eval form, then no question gets released. The current percentage of you who have filled in the course evals is 3% (so a bit to go). I will post weekly updates on the response rate.

Course Evaluations





28%



11 Days Remaining Ends 12/8/2016 at 11:59 PM EST Available 12/25/2016 at 11:59 PM EST

Response Rate Enrolled Students

CS Ed week (Dec 5)



We need volunteers!

with the Department of Computer Science and Engineering at UB:

Children K-12 are invited to:

We need demos!



Monday, Dec. 5 | Davis Hall, UB



When to use Dynamic Programming

There are polynomially many sub-problems

Richard Bellman

Optimal solution can be computed from solutions to sub-problems

$$OPT(j) = max \{ v_j + OPT(p(j)), OPT(j-1) \}$$

There is an ordering among sub-problem that allows for iterative solution

OPT (j) only depends on OPT(j-1), ..., OPT(1)

Scheduling to min idle cycles

n jobs, ith job takes wi cycles

You have W cycles on the cloud



What is the maximum number of jobs you can schedule?

Subset sum problem

Input: n integers $w_1, w_2, ..., w_n$

bound W

Output: subset S of [n] such that

(1) sum of w_i for all i in S is at most W

(2) w(S) is maximized

Recursive formula

```
OPT(j, W') = max value out of w<sub>1</sub>,...,w<sub>i</sub> with bound W'
If w_i > W'
   OPT(j, W') = OPT(j-1, W')
else
   OPT(j, W') = max \{ OPT(j-1, W'), w_i + OPT(j-1,W'-w_i) \}
```

Today's agenda

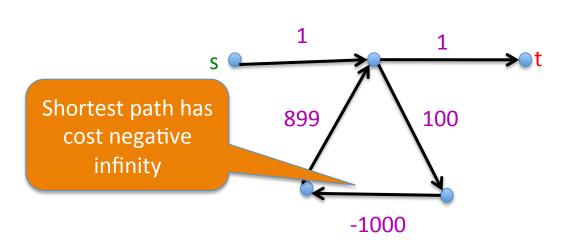
Dynamic Program for Subset Sum problem

Shortest Path Problem

Input: (Directed) Graph G=(V,E) and for every edge e has a cost c_e (can be <0)

t in V

Output: Shortest path from every s to t



Assume that G has no negative cycle

May the Bellman force be with you

