

Lecture 7

CSE 331

Sep 14, 2016

Check your PDF submits

Make sure to preview you PDF submission to Autolab!

A corrupted PDF file will get you a zero on that question

GS algo outputs a stable matching

Last lecture, GS outputs a perfect matching S

Lemma 3: S has no instability

Proof technique de jour

Proof by contradiction

Assume the negation of what you want to prove

After some
reasoning



Two observations

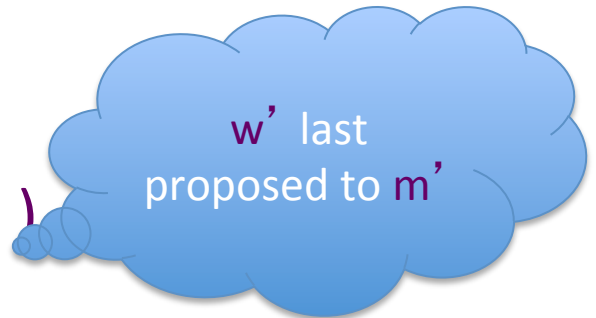
Obs 1: Once m is engaged he keeps getting engaged to “better” women

Obs 2: If w proposes to m' first and then to m (or never proposes to m) then she prefers m' to m

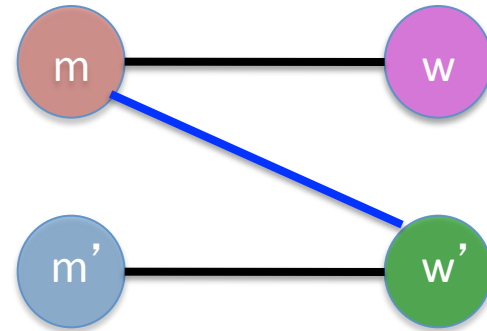
Proof of Lemma 3

By contradiction

Assume there is an instability (m, w')



m prefers w' to w
 w' prefers m to m'



Contradiction by Case Analysis

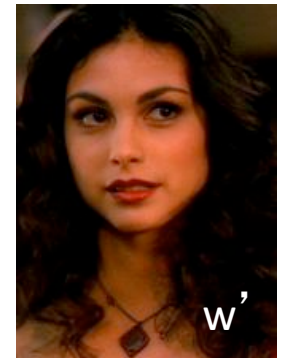
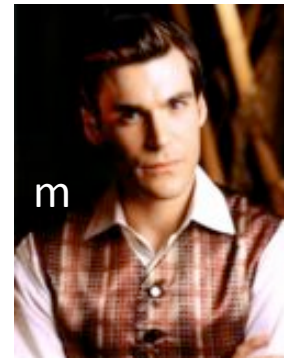
Depending on whether w' had proposed to m or not

Case 1: w' never proposed to m

w' prefers m' to m

By Obs 2

Assumed w' prefers m to m'



Case 2: w' had proposed to m

Case 2.1: m had accepted w' proposal

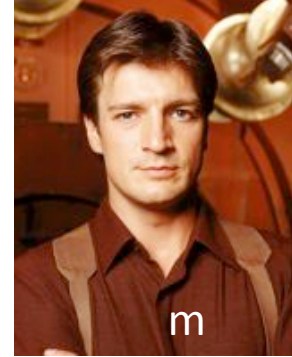
m is finally engaged to w

Thus, m prefers w to w'

By Obs 1



4simpsons.wordpress.com



Case 2.2: m had rejected w' proposal

m was engaged to w'' (prefers w'' to w')

By Obs 1

m is finally engaged to w (prefers w to w'')

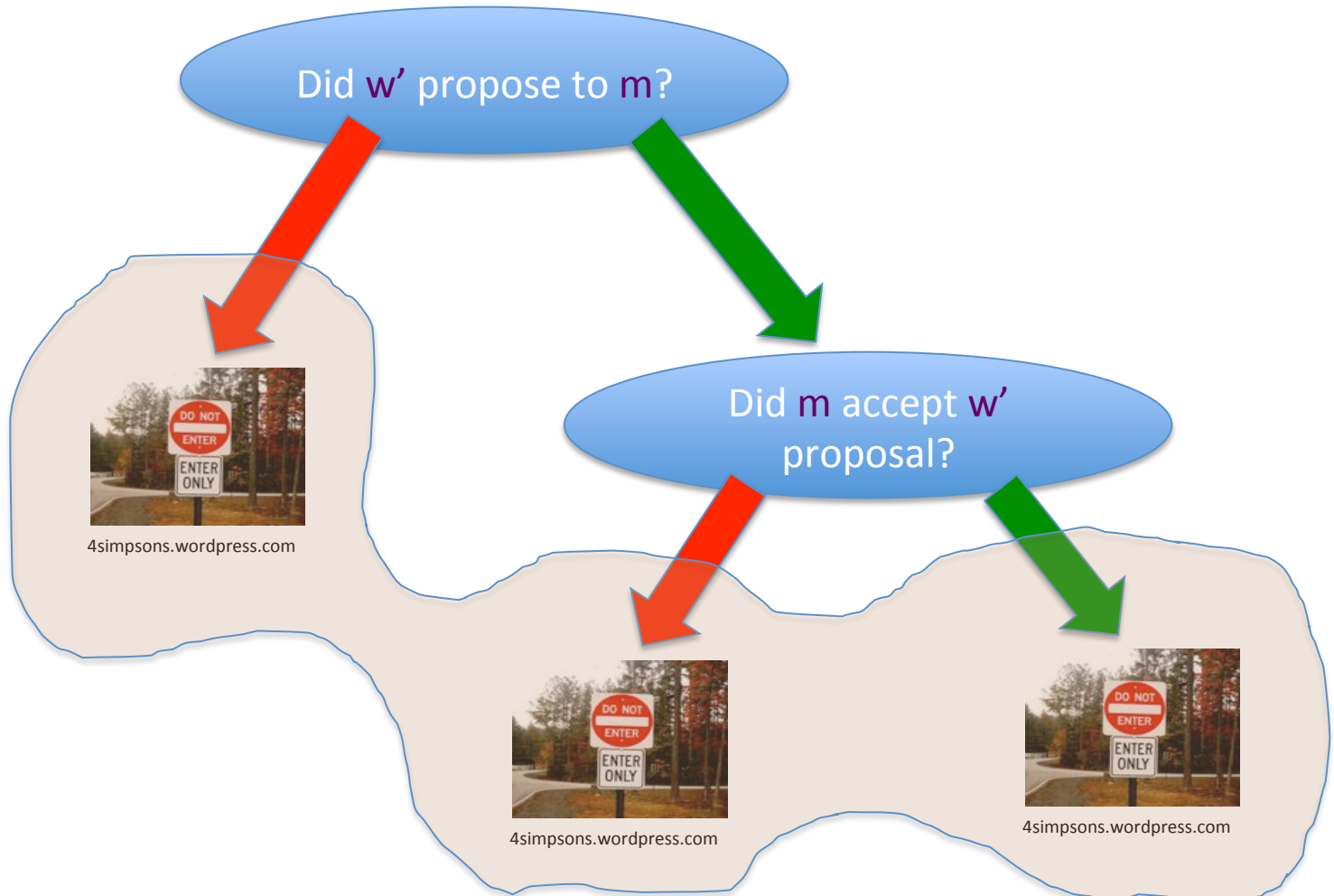
By Obs 1

m prefers w to w'



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Overall structure of case analysis



Questions?

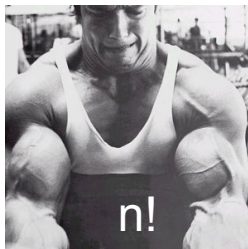
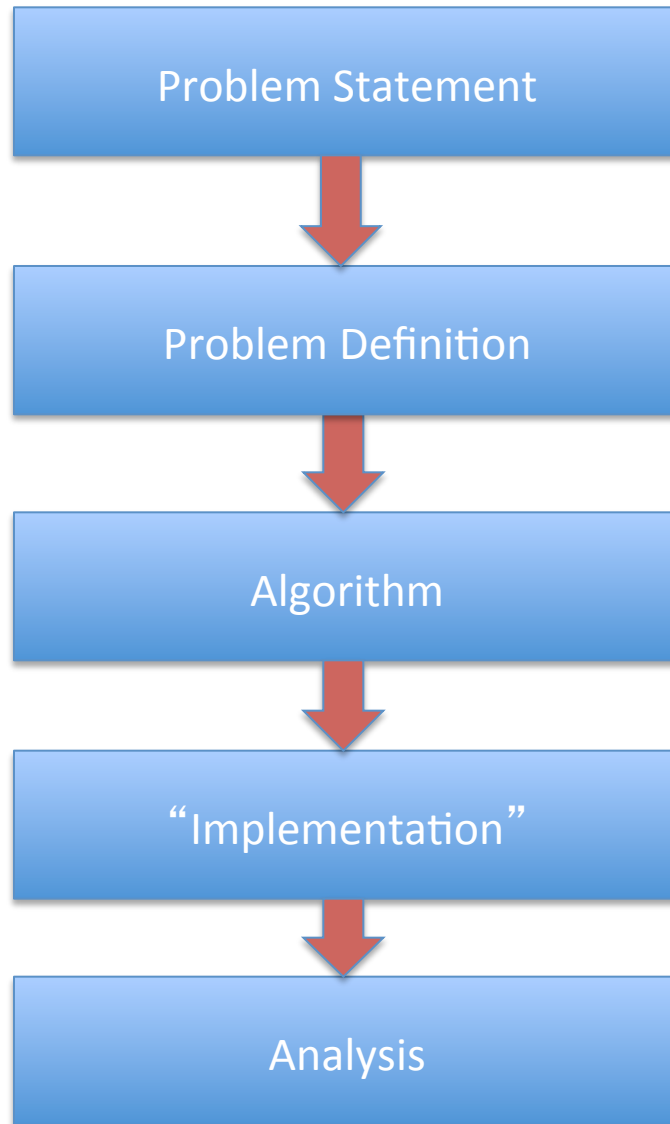


Extensions

Fairness of the GS algorithm

Different executions of the GS algorithm

Main Steps in Algorithm Design



Correctness Analysis

Definition of Efficiency

An algorithm is efficient if, when implemented, it runs quickly on real instances

Implemented where?



Platform independent definition

What are real instances?

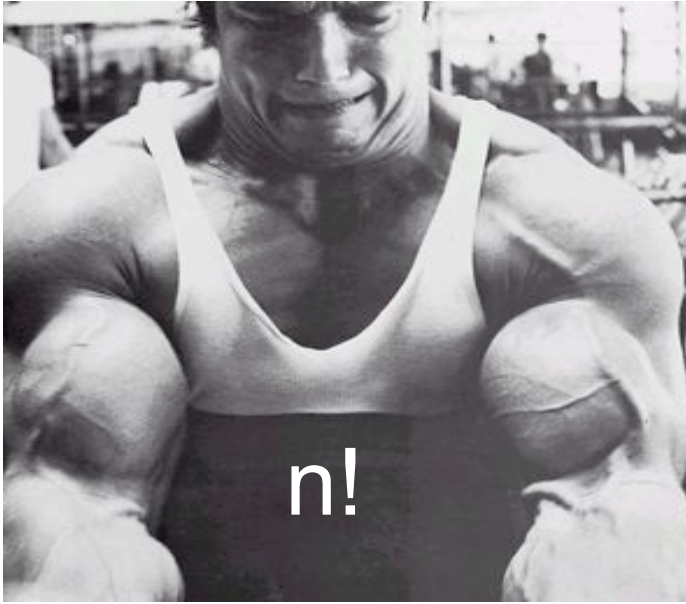
Worst-case Inputs

$$N = 2n^2 \text{ for SMP}$$

Efficient in terms of what?

Input size N

Definition-II



Analytically better than brute force

How much better? By a factor of 2?

Definition-III

Should scale with input size

If N increases by a constant factor,
so should the measure



Polynomial running time

At most $c \cdot N^d$ steps ($c > 0$, $d > 0$ absolute constants)

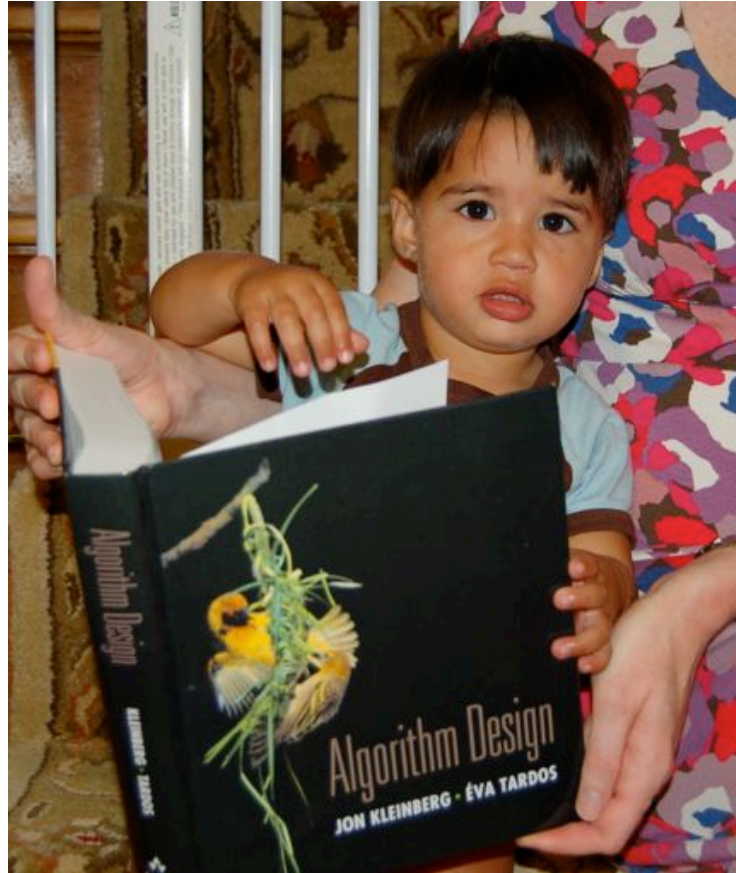
Step: “primitive computational step”

More on polynomial time

Problem centric tractability

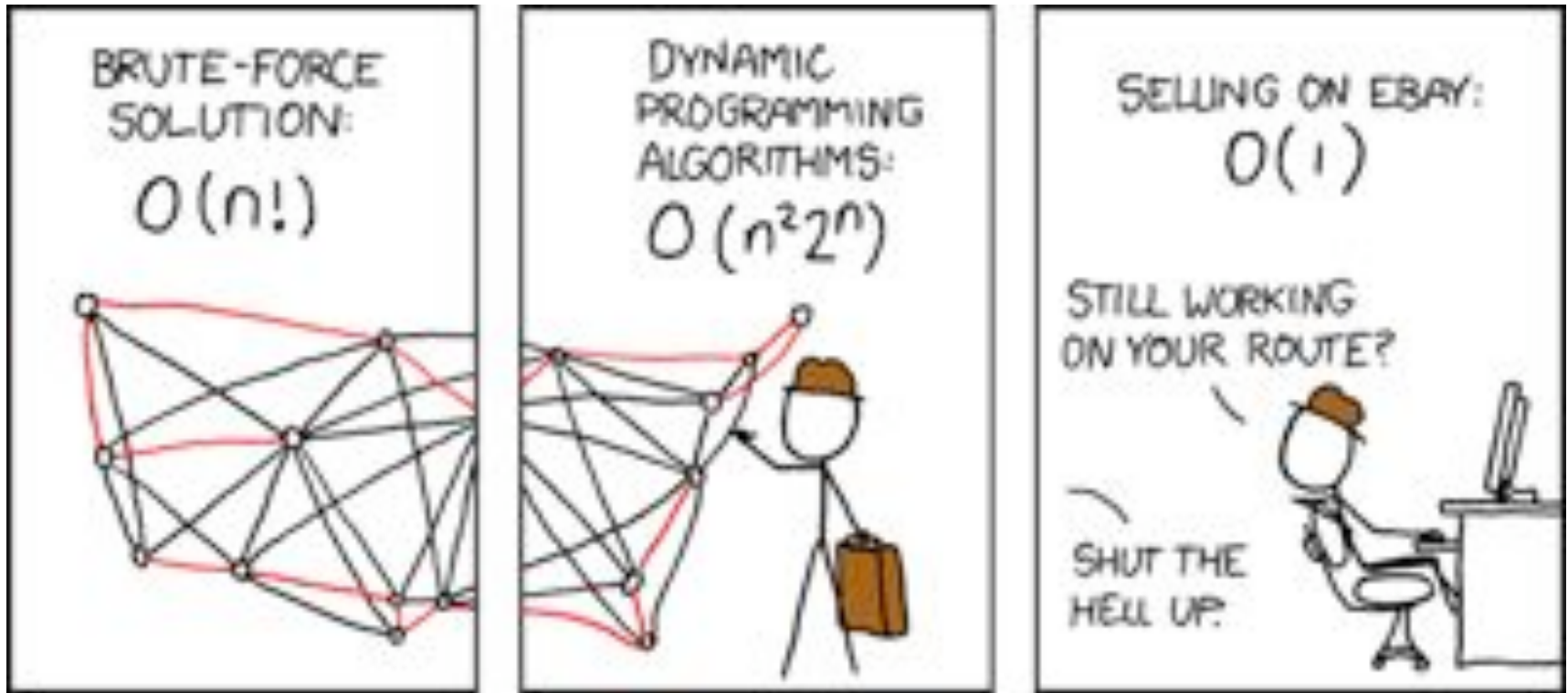
Can talk about problems that are not efficient!

Reading Assignments



Sections 1.2, 2.1, 2.2 and 2.4 in [KT]

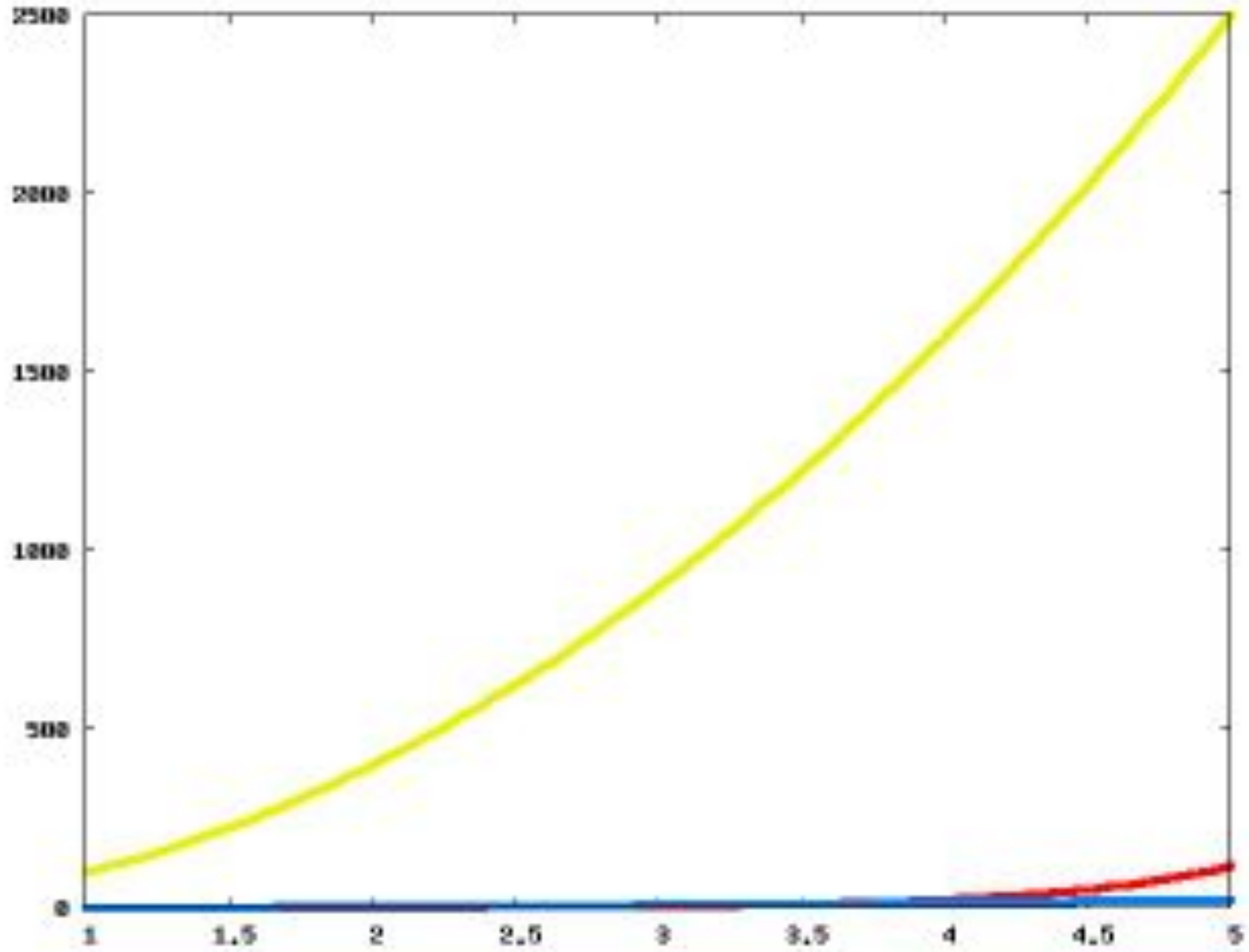
Asymptotic Analysis



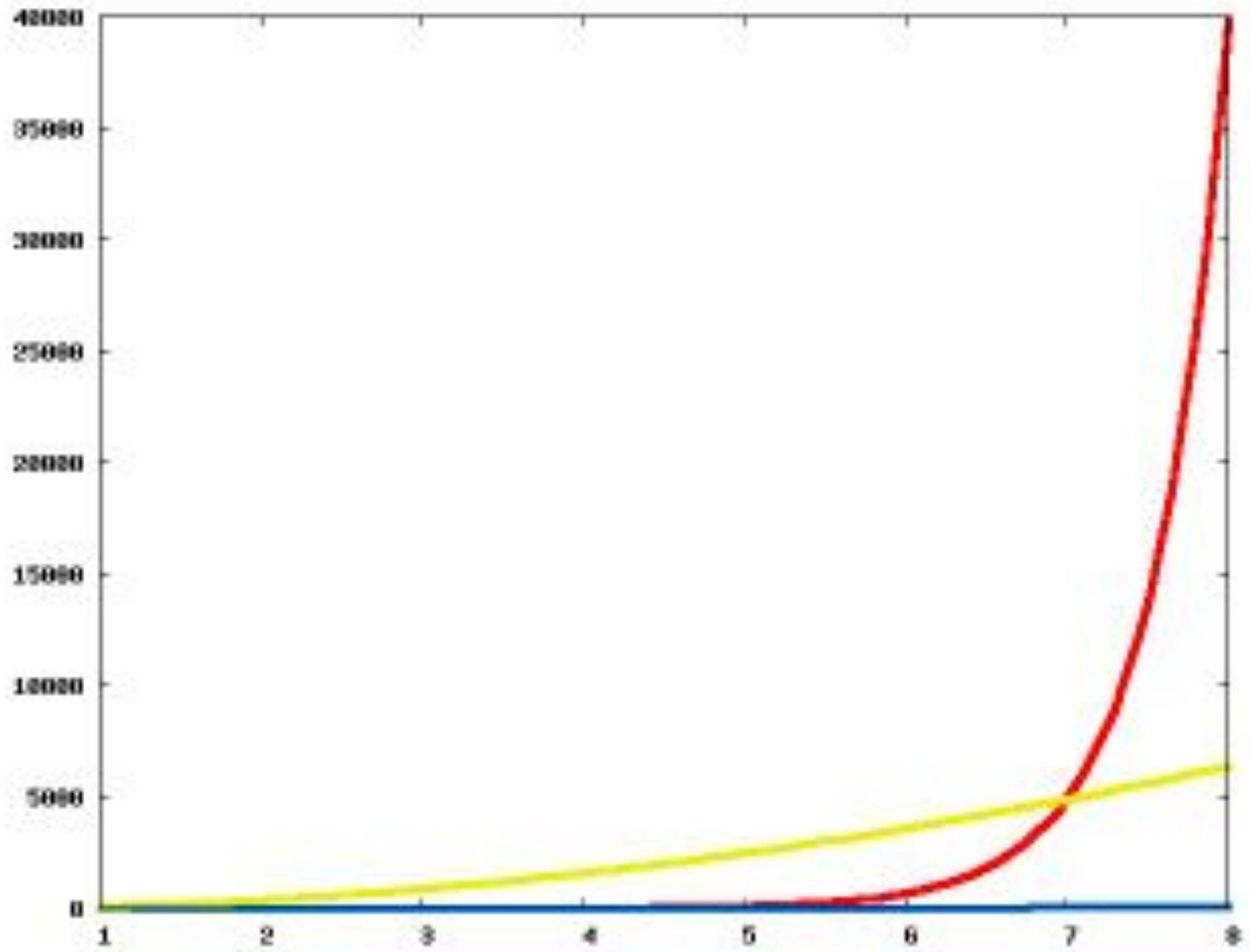
Travelling Salesman Problem

(<http://xkcd.com/399/>)

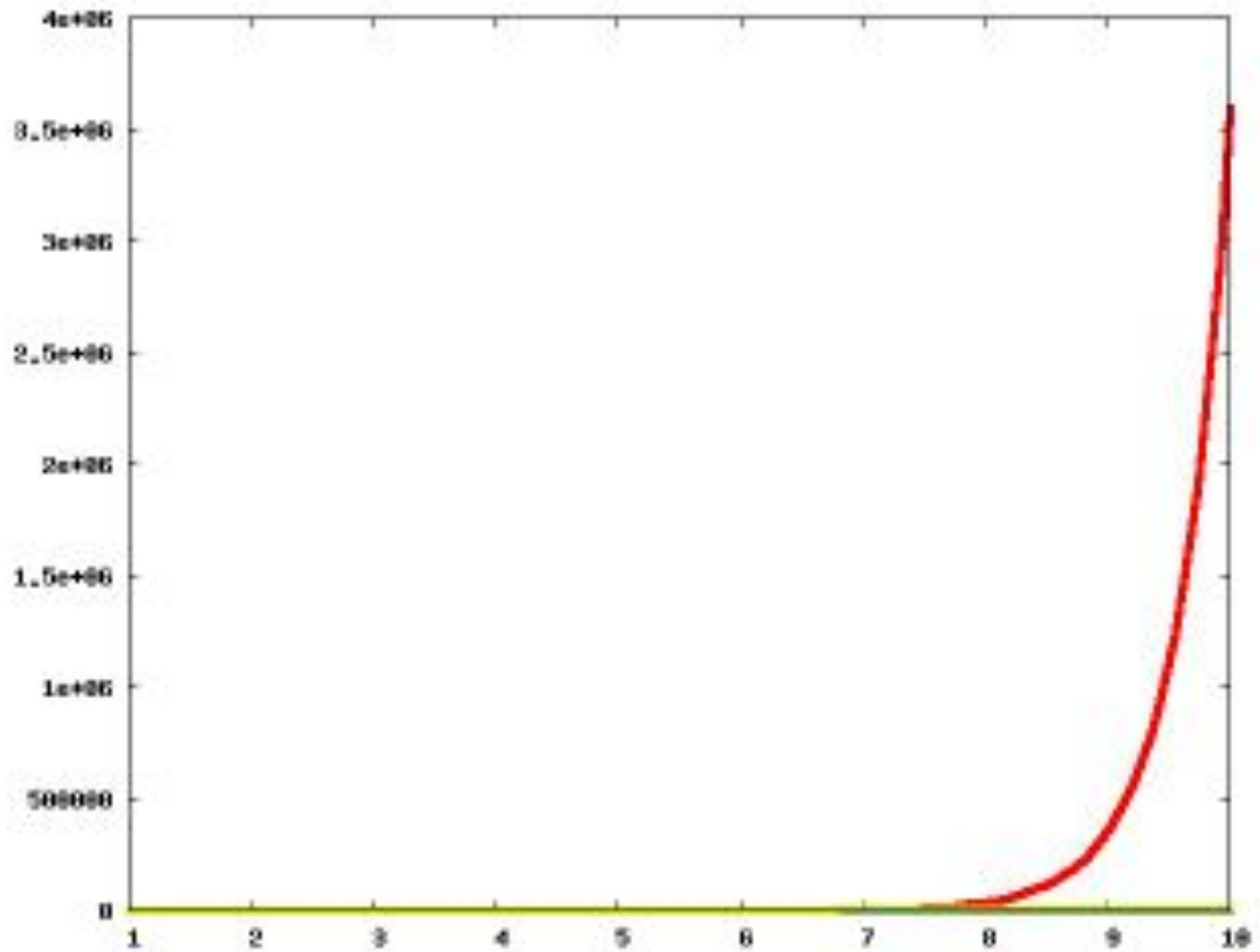
Which one is better?



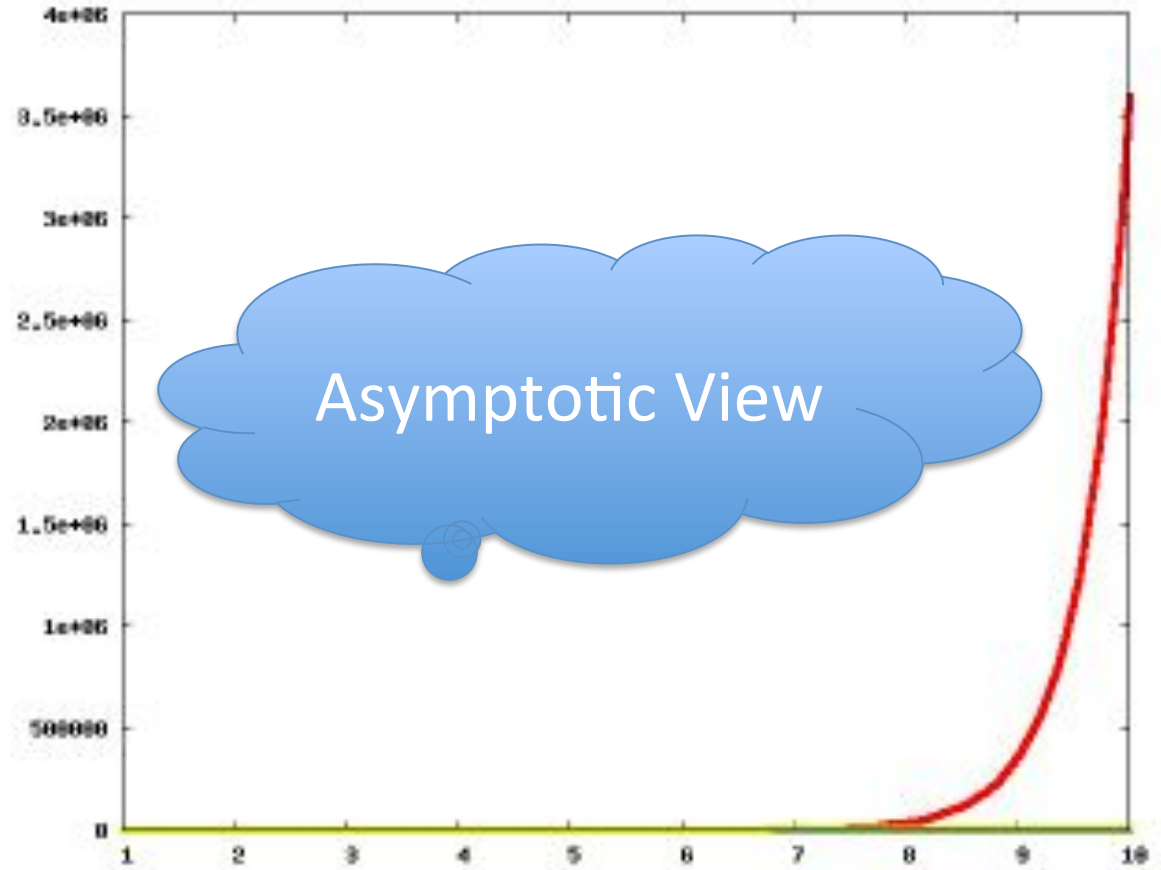
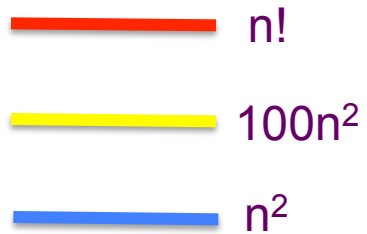
Now?



And now?



The actual run times



Asymptotic Notation



\leq is O with glasses

\geq is Ω with glasses

$=$ is Θ with glasses

Another view

remain anonymous on the web, let me know).

Silly way to remember
asymptotic notation....
Stick figure:



Big O

"^{head}ceiling of functn"

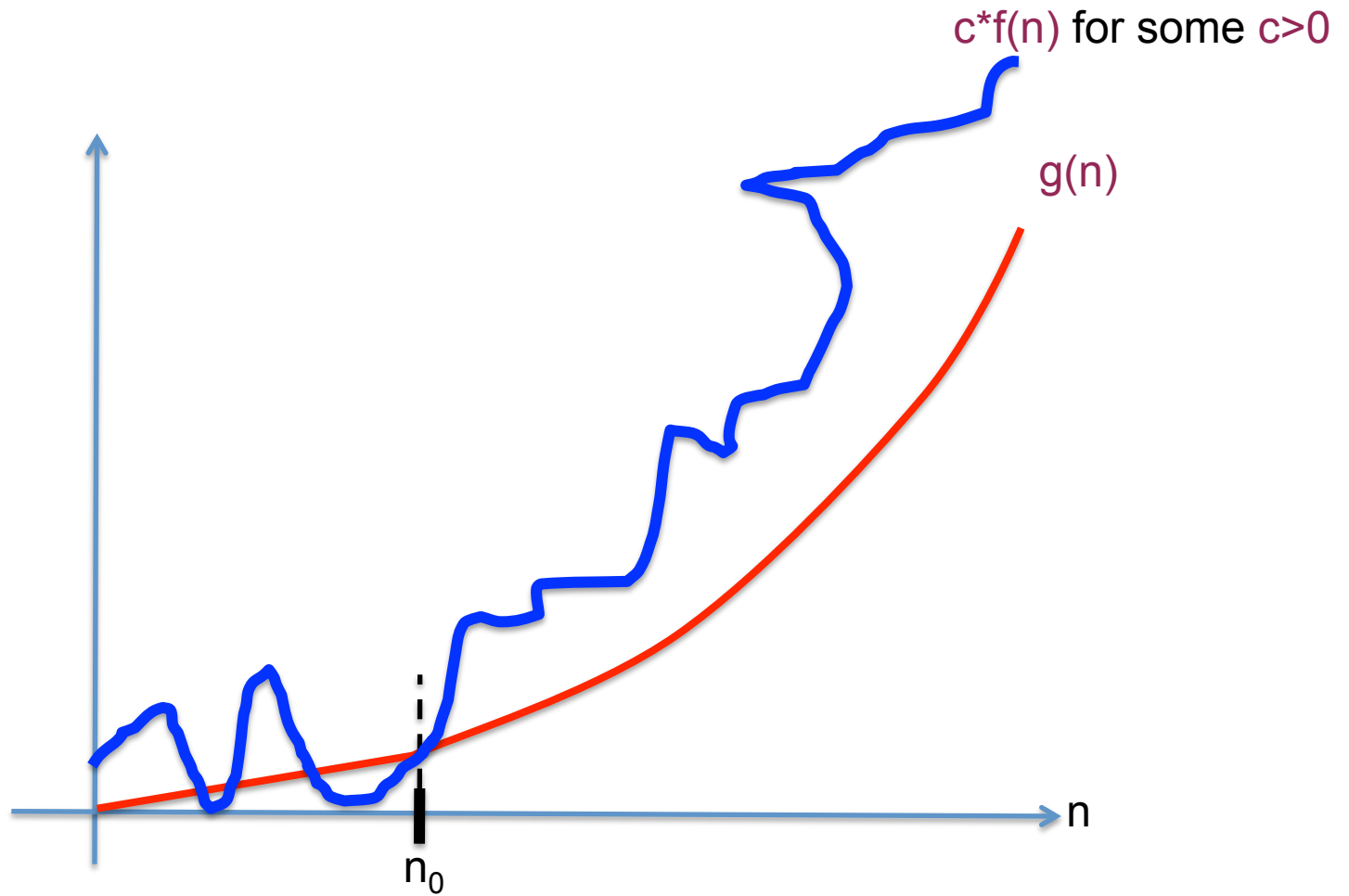
Big Θ

B/w Big- O + Big- Ω

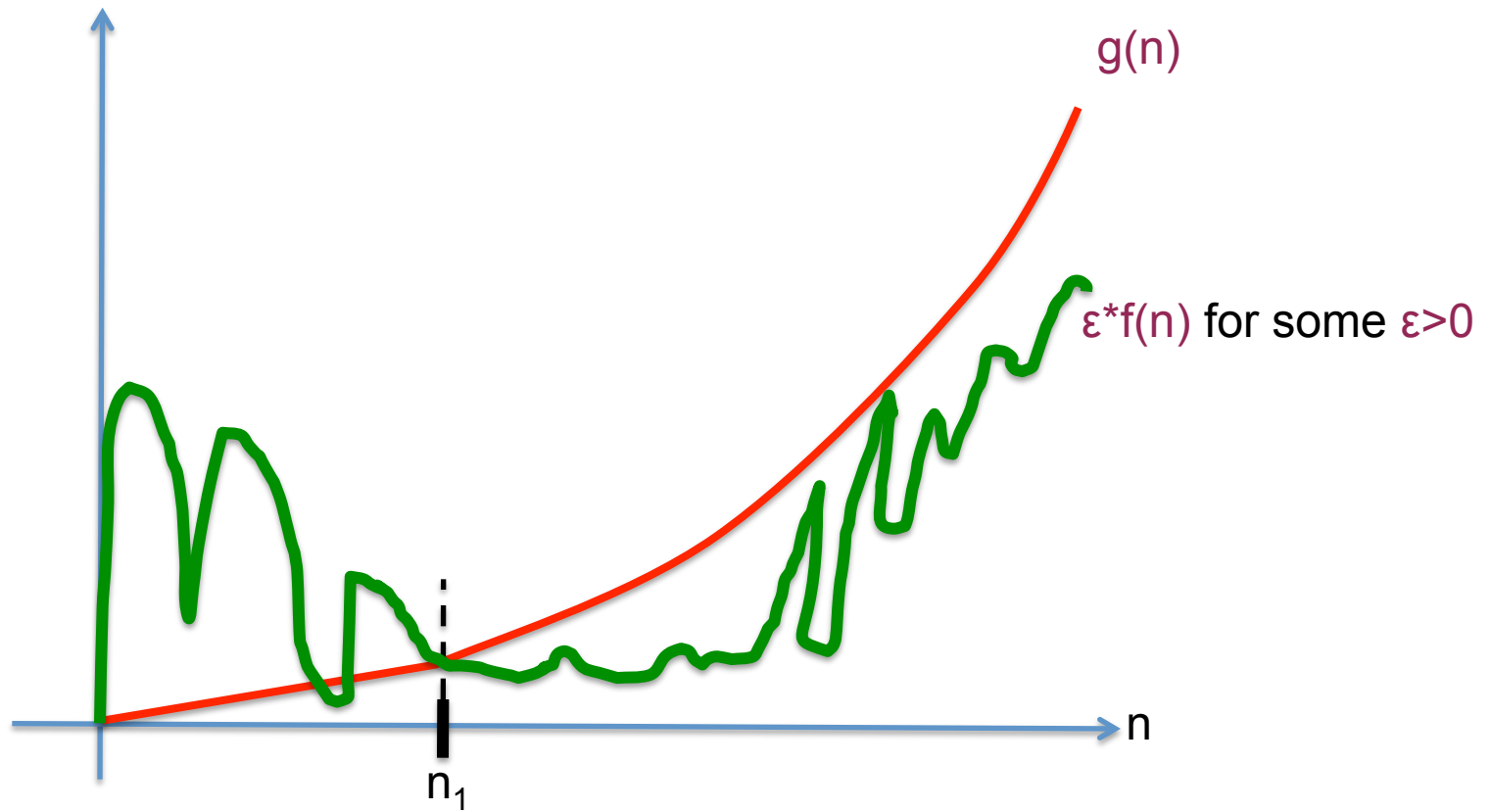
Big Ω

"Floor of functn"
feet

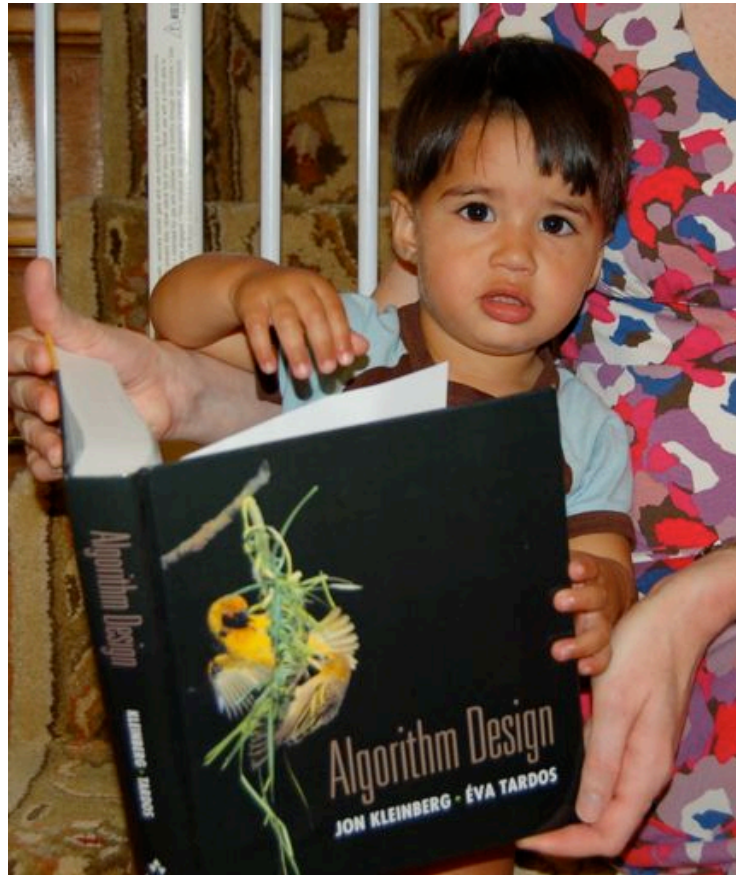
$g(n)$ is $O(f(n))$



$g(n)$ is $\Omega(f(n))$



Reading Assignments



Sections 1.1, 1.2, 2.1, 2.2 and 2.4 in [KT]