#### Lecture 27

CSE 331 Nov 3, 2017

# **UB Hacking**



# You need to re-form groups

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# HW 7 posted

#### Homework 7

Due by 11:00am, Friday, November 10, 2017.

Make sure you follow all the homework policies.

All submissions should be done via Autolab.

#### Question 1 (Programming Assignment) [40 points]

#### </> Note

This assignment can be solved in either Java, Python or C++ (you should pick the language you are most comfortable with). Please make sure to look at the supporting documentation and files for the language of your choosing.

#### The Problem

In this problem, we will explore weighted graphs.

We are given a starting node s and an ending node e, for some undirected graph G with n nodes. Further, each node u has its own weight,  $w_u$  (0 <=  $w_u$  <= 50). The graph

## Solutions for HW 6

#### At the END of the lecture

# Mergesort algorithm

Input: a<sub>1</sub>, a<sub>2</sub>, ..., a<sub>n</sub>

Output: Numbers in sorted order

MergeSort(a, n)

If n = 1 return the order  $a_1$ 

a<sub>L</sub> = a<sub>1</sub>,..., a<sub>n/2</sub>

 $a_{R} = a_{n/2+1}, ..., a_{n}$ 

return MERGE ( MergeSort(a<sub>L</sub>, n/2), MergeSort(a<sub>R</sub>, n/2) )

### Correctness

Input: a<sub>1</sub>, a<sub>2</sub>, ..., a<sub>n</sub>

Output: Numbers in sorted order



Inductive step follows from correctness of MERGE

## Rest of today's agenda

Analyze runtime of mergesort algorithm

# **Divide and Conquer**

Divide up the problem into at least two sub-problems

Recursively solve the sub-problems

"Patch up" the solutions to the sub-problems for the final solution

### Improvements on a smaller scale

Greedy algorithms: exponential  $\rightarrow$  poly time

(Typical) Divide and Conquer:  $O(n^2) \rightarrow$  asymptotically smaller running time

# Multiplying two numbers

Given two numbers a and b in binary

 $a=(a_{n-1},...,a_0)$  and  $b = (b_{n-1},...,b_0)$ 

Compute c = a x b

Running time of primary school algorithm?