Lecture 36

CSE 331 Nov 29, 2017

Quiz 2 on Monday



stop following





Quiz 2 on December 4

A gentle reminder that quiz 2 will be in class on Monday, December 4 from 1-1:10pm. (This is the Monday in last week of class.)

The first two questions will be T/F without justification (so like two from Q1 on sample final- 0842) and the third question will be T/F with justification (so like one from Q2 on sample final- 0842 but with the modification below).

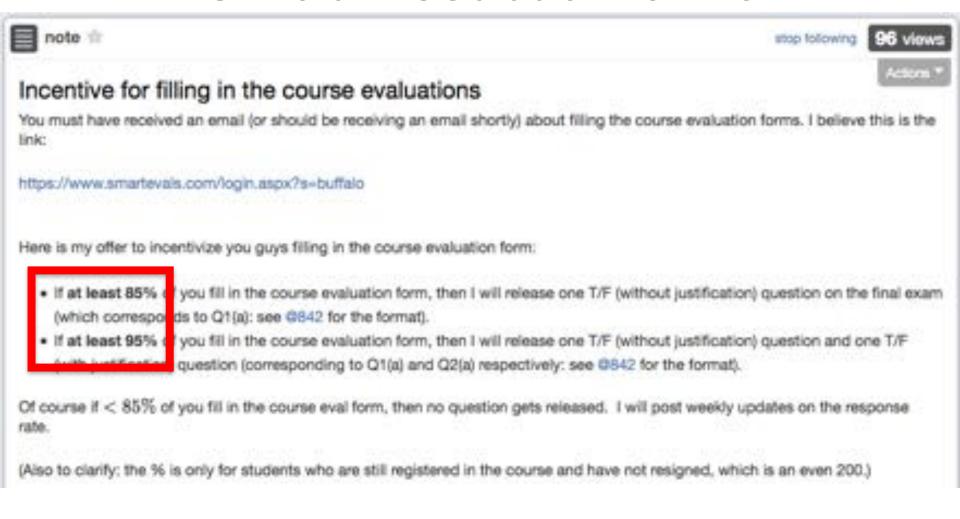
Based on the suggestion in @806, the T/F with justification question will be of the following format:

- You will be given a correct statement and will be asked to justify it (2 points)
- Then you will be given a variant of the correct statement and will be asked to say whether this statement if True or False and you will need to prove justification for your claim.
 - . Correct T/F will be worth 1 points and the justification will be worth 3 points.
 - Incorrect T/F will get 0 out of 4 irrespective of the justification.

You can bring in two 8.5° X 11° review sheets (you can use all four sides).

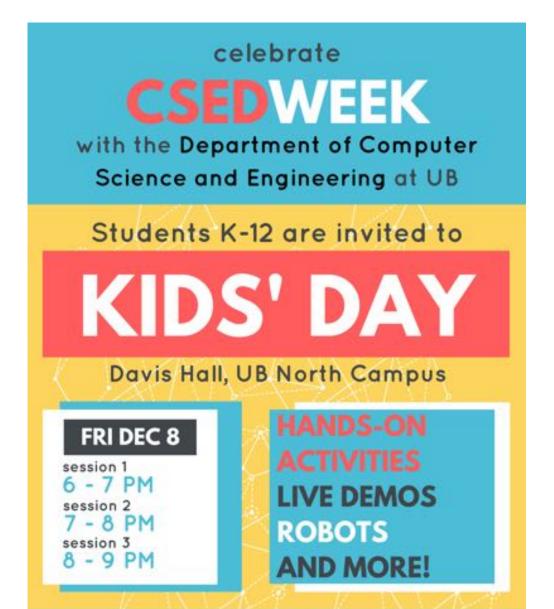
#pin

Official Feedback forms



FA 2017 2179_10751_Y CSE 331LR Intro to Algorithms LEC 200 48 152 24 % 11/24/17 12/10/1

We need demos for Nov 30!



Scheduling to min idle cycles

n jobs, ith job takes w_i cycles

You have W cycles on the cloud



What is the maximum number of jobs you can schedule?

When to use Dynamic Programming

O(nW) runtime

There are polynomially many sub-problems

$$OPT(j,B)$$
 $0 \le j \le n$, $0 \le B \le W$



Richard Bellman

Optimal solution can be computed from solutions to sub-problems

$$OPT(j, B) = ...$$

There is an ordering among sub-problem that allows for iterative solution

OPT (j,B) only depends on OPT(j-1, 0), ..., OPT(j-1,B)

Is O(nW) polynomial time?

n jobs, ith job takes wi cycles

NO. Pseudo-polynomial

You have W cycles on the cloud



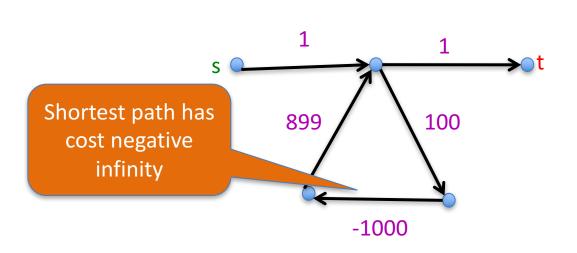
log W bits needed

What is the maximum number of jobs you can schedule?

Shortest Path Problem

Input: (Directed) Graph G=(V,E) and for every edge e has a cost c_e (can be <0)

Output: Shortest path from every s to t



Assume that G has no negative cycle

Today's agenda

Dynamic Program for shortest path

May the Bellman force be with you

