

Oct 8

THM 1: S^* is an optimal solution

↳ among all valid schedules with max # intervals

Ex 1: Algo terminates

Ex 2: S^* is a valid schedule

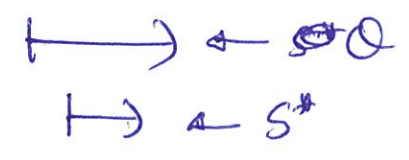
Pf of correctness of greedy algo
 ↳ Greedy stays ahead (next)
 ↳ Exchange argument (later)

Let Θ be an optimal solution

Ex: Convince yourself that such an optimal solution \exists .

Idea 1: $S^* = \Theta$ } problem: have multiple optimal solutions

Idea 2: Show $|S^*| = |\Theta|$



THM 2: $|S^*| = |\Theta|$

Notation: $S^* = \{i_1, \dots, i_k\}$

$f(i_1) \leq f(i_2) \leq \dots \leq f(i_k)$

$\Theta = \{j_1, \dots, j_m\}$

$f(j_1) \leq \dots \leq f(j_m)$

THM 2': $k = m$

Claim: $k \leq m$

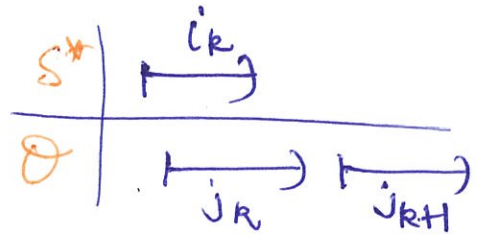
Lemma 1 ("Greedy stays ahead") $\forall 1 \leq l \leq k$

$f(i_l) \leq f(j_l)$

Pf idea for Thm 2': (Assume Lemma 1.)

Pf by contradiction: Assume $k \neq m$

\Rightarrow Claim $k < m \Leftrightarrow m \geq k+1$
 $\Rightarrow j_{k+1} \in \mathcal{Q}$



By Lemma 1, $f(i_k) \leq f(j_k)$

Consider the case when Greedy adds

i_k to S

$j_{k+1} \in R \rightarrow R \neq \emptyset \Rightarrow$ Greedy did not terminate
as j_{k+1} does not conflict with anything in $S^* \Rightarrow$ contradiction