# Lecture 17 

CSE 331
Oct 4, 2019

## HW 5 is out

## Homework 5

Due by 11:00am, Friday, October 11, 2019.
Make sure you follow all the homework policies.
All submissions should be done via Autolab.
The care package on minimizing the maximum lateness problem would be useful for Q3 and might be useful for Q2(b) as well.

## Question 1 (High Speed Internet) [50 points]

## The Problem

We come back to the issue of many USA regions not having high speed internet. In this question, you will consider an algorithmic problem that you would need to solve to help out a (fictional) place get high speed Internet.

You are the algorithms whiz in the effort to bring high speed Internet to SomePlaceInUSA. After lots of rounds of discussions and public feedback, it was decided that the most cost-effective way to bring high speed internet to SomePlaceInUSA was to install high speed cell towers to connect all houses in SomePlaceInUSA to high speed internet. There are two things in your favor:

1. It just so happens that all of the $n$ houses in SomePlaceInUSA are on the side of a straight road that runs through the town.
2. The above implies that you only need cell towers that only need to broadcast their signal in a narrow range, which means one cell tower can provide high speed internet
[^0]
## HW 4 Solutions

At the end of the lecture

## Graded HW 3

## Perhaps by tonight?

## Quiz on Monday

## Quiz 1 on Monday, Oct 7

The first quiz will be from 1-1:10pm in class on Monday, October 7 . We will have a 5 mins break after the quiz and the lecture will start at $1: 15 \mathrm{pm}$.

We will hand out the quiz paper at 12:55pm but you will NOT be allowed to open the quiz to see the actual questions till 1 pm. However, you can use those 5 minutes to go over the instructions and get yourself in the zone.

There will be two T/F with justification questions (like those in the sample mid term 1: @641.) Also quiz 1 will cover all topics we cover in class till Friday, Oct 4.

Also like the mid-term y'all can bring in one letter sized cheat-sheet (you can use both sides).
\#pin

## Update on coding project tonight

## Coding Mini Project

Problem 1 due at 11am, Friday, October 25, 2019.
Problems 2 and 3 due at 11am, Friday, November 22, 2019.
Problems 4 and 5 due at 11am, Friday, December 6, 2019.
All submissions should be done via Autolab.Acknowledgment
The development of the coding component of the mini-project was supported by a Mozilla Responsible Computer Science award [']. The support is gratefully acknowledged.

## Some Suggestions and Warnings

While this coding mini-project is somewhat similar to Question 3s on the homework, there are some crucial differences and we wanted to highlight few things for y'all upfront:

Form groups of size $\leq 3$
This is a group project (unlike Q3s on the HWs that had to be done individually) and you can work in groups of size at most 3 . The submissions will be on Autolab and everyone in the group will get the same grade. The project will be challenging so we highly recommend that you form a group of size at least 2 to make the workload reasonable.

## The "real" end of Semester blues



Write up a term paper

## Party!

Exam study
331 HW


## The "real" end of Semester blues



## Write up a term paper


Exam study


331 HW


## The algorithmic task



Write up a term paper


Project


## Scheduling to minimize lateness



## Write up a term paper


Exam study


## One possible schedule



## Minimizing Max Lateness

## Minimizing Maximum Lateness

This page collects material from previous incarnations of CSE 331 on scheduling to minimize maximum lateness.

## Where does the textbook talk about this?

Section 4.2 in the textbook has the lowdown on the problem of scheduling to minimize maximum lateness.

## Fall 2018 material

First lecture
Here is the lecture video:


## Today



## Reading Assignment

Sec 2.5 of [KT]


## Shortest Path problem

Input: Directed graph G=(V,E)
Edge lengths, $\mathrm{I}_{\mathrm{e}}$ for e in E

"start" vertex s in V


Output: All shortest paths from s to all nodes in $V$

## Naïve Algorithm

$\Omega(n!)$ time

## Dijkstra's shortest path algorithm




[^0]:    

