Lecture 18

CSE 331 Oct 15, 2021

Please have a face mask on

Masking requirement



<u>LIR_requires</u> all students, employees and visitors – regardless of their vaccination status – to wear face coverings while inside campus buildings.

https://www.buffalo.edu/coronavirus/health-and-safety/health-safety-guidelines.html

HW 4 is out

Homework 4

Due by 8:00am, Wednesday, October 20, 2021.

Make sure you follow all the homework policies.

All submissions should be done via Autolab.

The care package on minimizing the maximum lateness problem would be useful for Q3 and might be useful for Q2(b) as well.

Question 1 (High Speed Internet) [50 points)

Probably the easiest HW

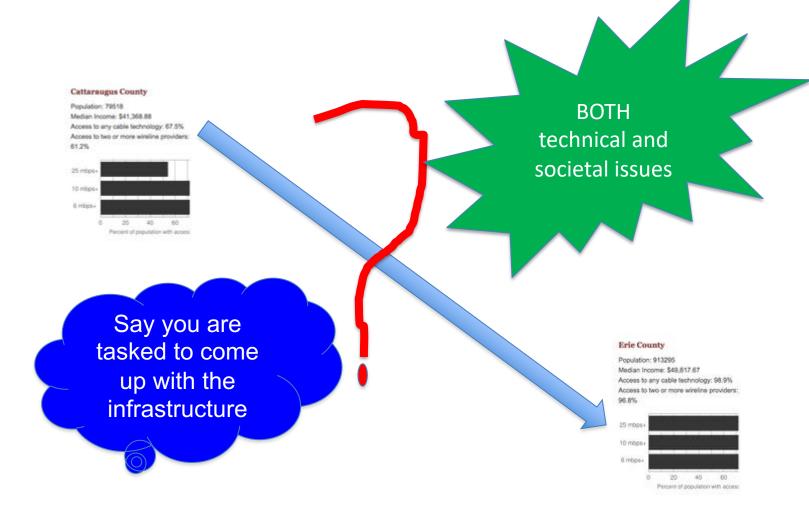
The Problem

We come back to the issue of many USA regions not having high speed internet. In this question, you will consider an algorithm problem this you would need to solve to help out a (fictional) place get high speed internet.

You are the algorithms whiz in the effort to bring high speed internet to SomePlaceInUSA. After lots of rounds of discussions and public feedback, it was decided that the most cost-effective way to bring high speed internet to SomePlaceInUSA was to install high speed cell towers to connect all houses in SomePlaceInUSA to high speed internet. There are two things in your favor:

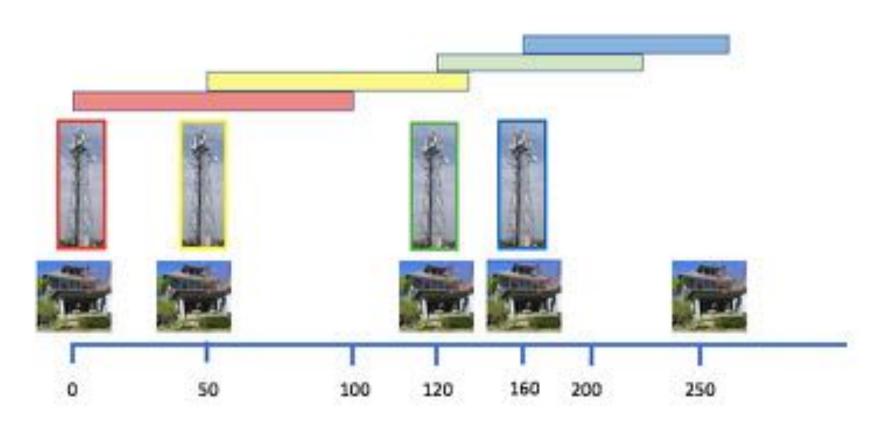
1. It just so happens that all of the n houses in SomePlaceInUSA are on the side of a straight road that runs through the town.

Make broadband more available



HW 4 Q1: How to lay down towers

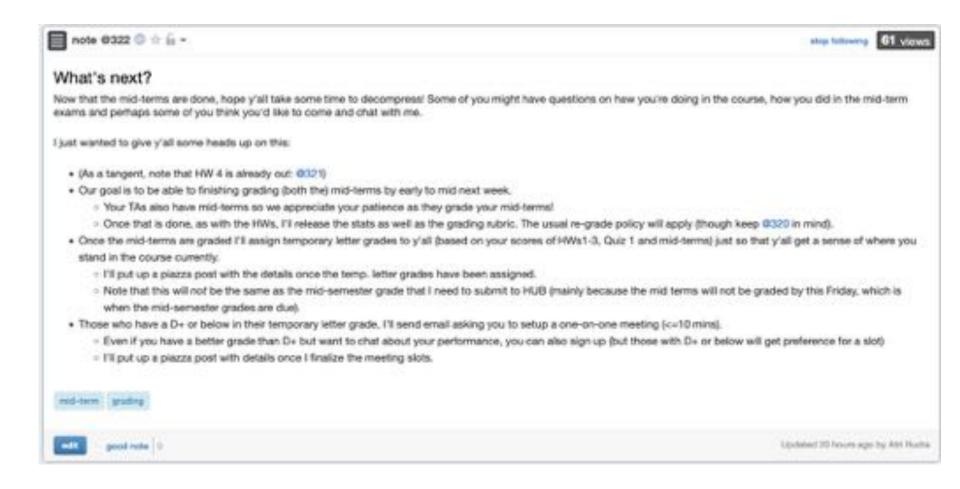
Here is a quick visual argument for the above leads to continuous cell coverage:



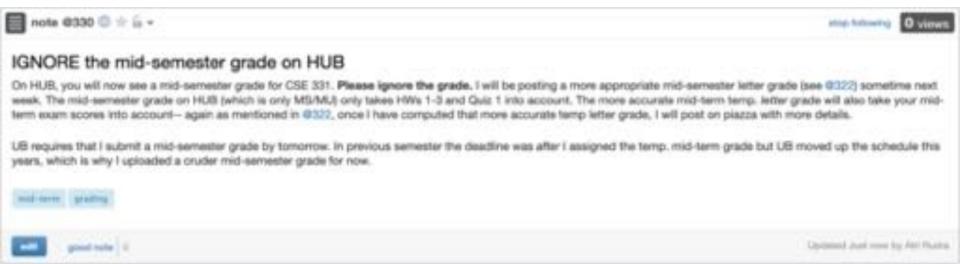
Project deadlines coming up

Fri, Oct 29	Counting Inversions P19 P18 P17 x3	[KT, Sec 5.3] (Project (Problem 1 Coding) in)
Mon, Nov 1	Multiplying large integers □F18 □F17 x²	[KT, Sec 5.5] (Project (Problem 1 Reflection) in) Reading Assignment: Unraveling the mystery behind the identity
Wed, Nov 3	Closest Pair of Points □F19 □F18 □F17 x3	[KT, Sec 5.4]
Fri, Nov 5	Kickass Property Lemma □F19 □F18 □F17 x²	[KT, Sec 5.4] (Project (Problem 2 Coding) in)
Mon, Nov 8	Weighted Interval Scheduling F19 F17 x²	[KT, Sec 6.1] (Project (Problem 2 Reflection) in)

Some other stuff coming up



Mid-semester grade on HUB



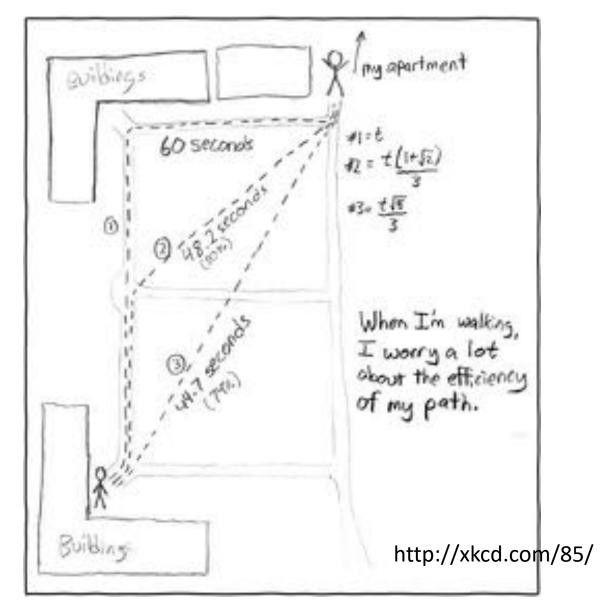
Note on re-grade requests



Questions?



Shortest Path Problem



Another more important application

Is BGP a known acronym for you?



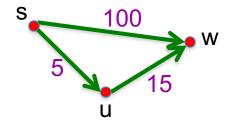
Routing uses shortest path algorithm

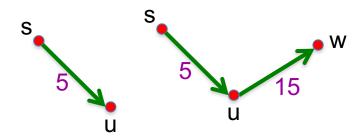
Shortest Path problem

Input: *Directed* graph G=(V,E)

Edge lengths, le for e in E

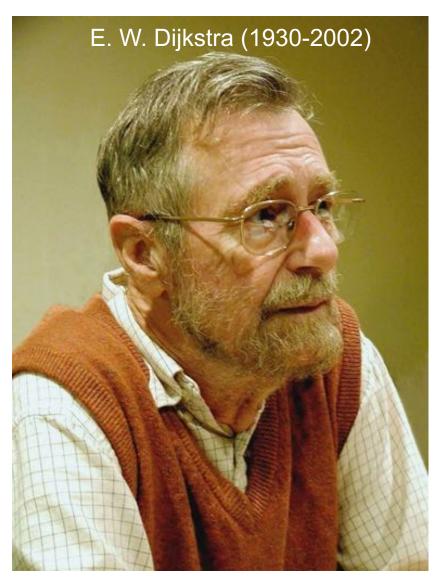
"start" vertex s in V



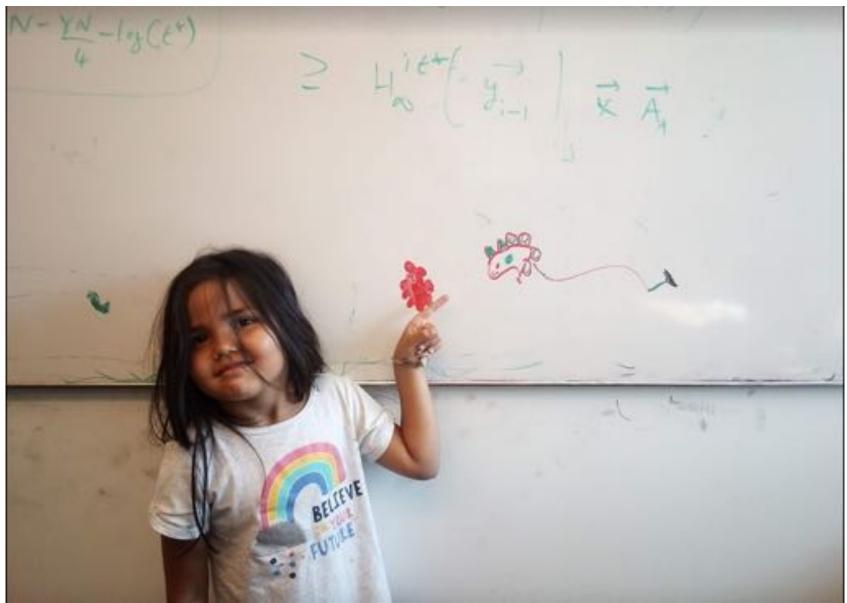


Output: Length of shortest paths from s to all nodes in V

Dijkstra's shortest path algorithm



Questions/Comments?



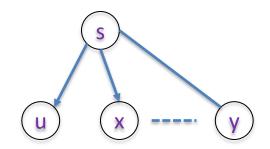
On to the board...



Towards Dijkstra's algo: part ek

Determine d(t) one by one

$$d(s) = 0$$



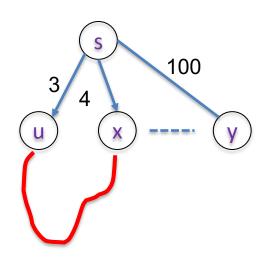
Towards Dijkstra's algo: part do

Determine d(t) one by one

Let u be a neighbor of s with smallest $I_{(s,u)}$

$$d(u) = I_{(s,u)}$$

Not making any claim on other vertices



Length of is
$$\ge 0$$

Towards Dijkstra's algo: part teen

Determine d(t) one by one

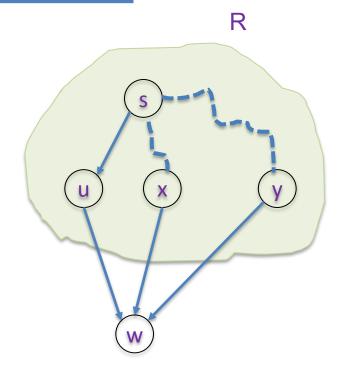
Assume we know d(v) for every v in R

Compute an upper bound d'(w) for every w not in R

$$d(w) \leq d(u) + I_{(u,w)}$$

$$d(w) \leq d(x) + I_{(x,w)}$$

$$d(w) \leq d(y) + I_{(y,w)}$$



$$d'(w) = \min_{e=(u,w) \text{ in E, u in R}} d(u)+I_e$$