

Oct 4

# Interval Scheduling Problem

$[3, 7)$   
 $= \{3, 4, 5, 6\}$

$f(i) - 1$  is finish time

Input:  $n$  intervals:  $i$ th interval  $[s(i), f(i))$

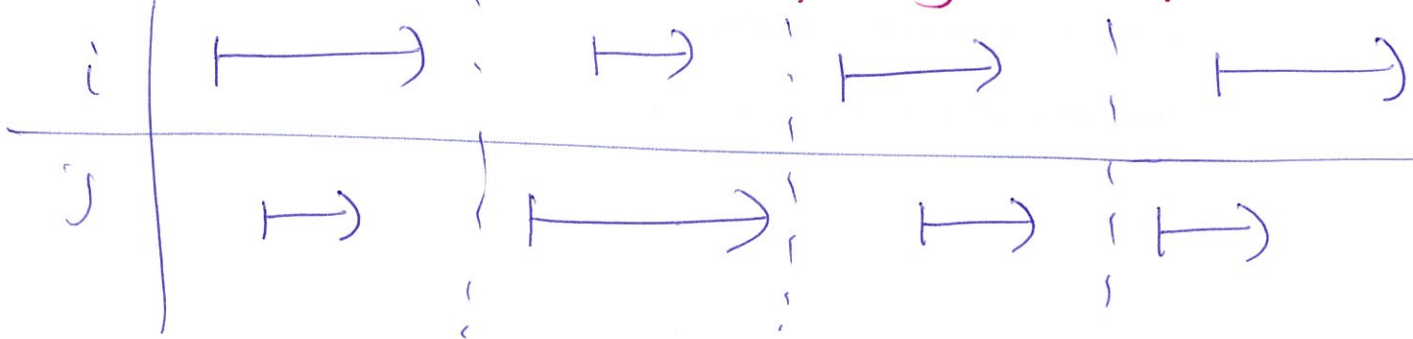
start time

Output: A valid schedule with max # intervals in

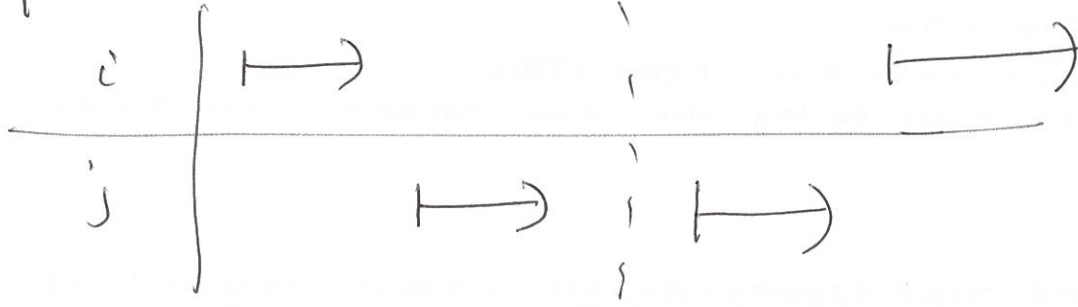
Def: A schedule  $S \subseteq [n] \stackrel{\text{def}}{=} \{i_1, \dots, i_k\}$  overall valid schedule

Def: A valid schedule  $S$  has no conflicts

Def: interval  $i$  &  $j$  conflict if they overlap



$\Rightarrow$  no conflict



Obs: A valid schedule sorted by start / finish time gives the same order.  $\rightarrow$  increasing order

Assume: Input intervals are sorted by finish time  
 $\Rightarrow f(1) \leq f(2) \leq f(3) \leq \dots \leq f(n)$

$\uparrow$  If not, sort in  $O(n \log n)$  time

## Greedy Algo

0.  $R \leftarrow [n]$

1.  $S \leftarrow \phi$

2. While  $R \neq \phi$

(2.1) Let  $i$  be the smallest index in  $R$

(2.2) Add  $i$  to  $S$

(2.3) Remove  $i$  from  $R$

(2.4) Delete all  $j \in R$  that conflict with  $i$

3. Return  $S^* \leftarrow S$

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THM:  $S^*$  is an optimal solution.

$\hookrightarrow$  i.e.  $\forall$  inputs, among all possible valid schedules for the input,  $S^*$  has the max. # of intervals.

Ex 1: Algo terminates

Ex 2:  $S^*$  is a ~~valid~~ valid schedule

Pf. of correctness of greedy algo  $\left\{ \begin{array}{l} \rightarrow \text{Greedy stays ahead (next)} \\ \rightarrow \text{Exchange argument (min. max. lateness)} \end{array} \right.$  (Sec 4.2)

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