

Oct 18

Lemma 1: At the end of each iteration,

if $u \in R$, the path P_u is a shortest $s-u$ path.
 $s-u$ path in Dijkstra tree

Pf(Idea): By induction on $|R|$

Base case: $|R|=1 \Rightarrow R = \{s\}$, $d(s)=0$ ✓

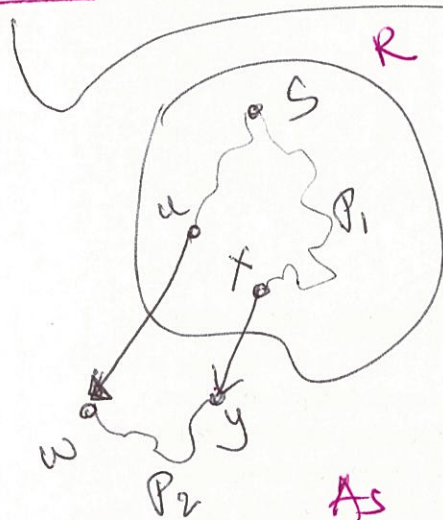
I.H.: Assume lemma is true for $|R|=k$

$(k \geq 1)$

I.S.: Assume $|R|=k+1$

Let w be the $(k+1)^{th}$ vertex added to R

Goal: P_w is a shortest $s-w$ path. assume w was discovered from u



$P_w = P_u, w$ our construction of Dijkstra tree + $\{u,w\}$

Pf(idea) By contradiction.

Assume \exists an $s-w$ path $P'_w \neq P_w$ but $l(P'_w) < l(P_w)$ (*)

As $s \in R$ but $w \notin R \Rightarrow P'_w$ "crosses" R at some edge (x,y)

$P'_w = P_1, x, y, P_2$

$$d'(w) = \min_{\substack{(u,w) \in E \\ u \in R}} \left\{ \begin{array}{l} d(u) + l_{(u,w)} \\ l(P_2) \geq 0 \end{array} \right.$$

$$l(P'_w) = l(P_1) + l_{(x,y)} + l(P_2) \stackrel{\text{def. of } d'(y)}{\geq} d(x) + l_{(x,y)} + l(P_2) \geq d'(y) + l(P_2) \geq d'(y)$$

$d(x)$ is distance from s to x

$$\geq d'(w)$$

Dijkstra chose w over y Algo def $\Rightarrow d'(w) = l(P_w)$

$\Rightarrow l(P'_w) \geq l(P_w) \Rightarrow$ contradiction \square