

Oct 4

# Interval Scheduling Problem

[3,8) = {3, ~~4~~, 5, 6, 7}

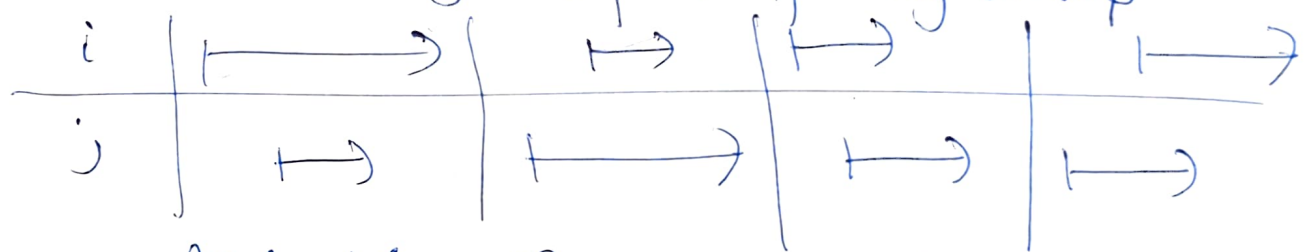
Input:  $n$  intervals:  $i^{\text{th}}$  interval  $[s(i), f(i))$   
 $1 \leq i \leq n$

Output: A valid schedule with the max number of intervals in it.  
 $= \{s(i), \dots, f(i)-1\}$

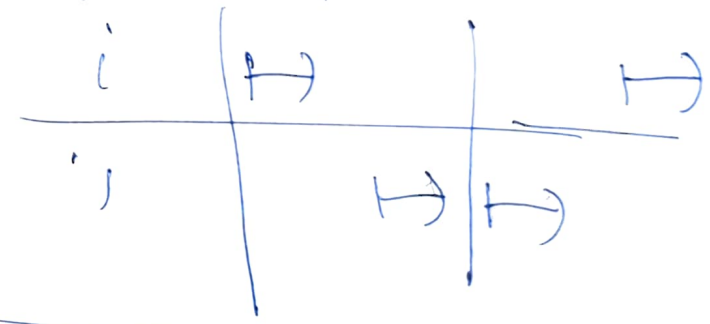
Def: A schedule  $S \subseteq [n]$  def  $\{1, \dots, n\}$

Def: A valid schedule  $S$  has no conflicts

Def: Intervals  $i$  &  $j$  conflict if they overlap



$\Rightarrow$  no conflict b/w  $i$  &  $j$



Claim: A valid schedule that is sorted by start time or by finish time gives the same order



Assume: Input intervals are sorted by their finish times  
 $f(1) \leq f(2) \leq f(3) \leq \dots \leq f(n)$

$\hookrightarrow$  If not: in  $O(n \log n)$  time sort the intervals.

# Greedy algo

0.  $R \leftarrow [n]$

1.  $S \leftarrow \emptyset$

2. While  $R \neq \emptyset$

(2.1) Let  $i$  be the smallest index in  $R$

(2.2) Add  $i$  to  $S$

*Combine* { ~~(2.3)~~ Remove  $i$  from  $R$

(2.4) Delete all  $j \in R$  (from  $R$ ) that conflict with  $i$

3. Return  $S^* \leftarrow S$

OK since  $f(i) \leq f(j) \leq \dots \leq f(n)$

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Theorem 1:  $S^*$  is an optimal solution  $\rightarrow$   $\forall$  inputs, among all possible valid schedules for that input,  $S^*$  has the max number of intervals.

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Pf. of correctness  $\left\{ \begin{array}{l} \rightarrow \text{Greedy stays ahead (next)} \\ \rightarrow \text{Exchange argument (min. max lateness} \rightarrow \text{Sec 4.2)} \end{array} \right.$